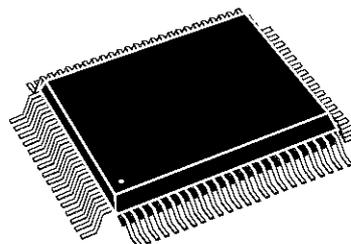
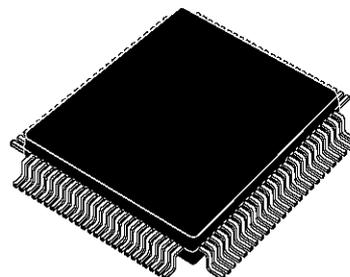


**16-BIT ROMLESS MICROCONTROLLER****ADVANCE DATA**

- High performance 16-bit CPU with 4-stage pipeline
- 100 ns instruction cycle time at 20-MHz CPU clock
- 500 ns multiplication (16 × 16 bits), 1 μs division (32 / 16 bit)
- Enhanced boolean bit manipulation facilities
- Additional instructions to support HLL and operating systems
- Register-based design with multiple variable register banks
- Single-cycle context switching support
- Up to 16 Mbytes linear address space for code and data
- 2 Kbytes on-chip RAM
- Programmable external bus characteristics for different address ranges
- 8-Bit or 16-bit external data bus
- Multiplexed or demultiplexed external address/data buses
- Five programmable chip-select signals
- Hold- and hold-acknowledge bus arbitration support
- 1024 bytes on-Chip special function register area
- Idle and power down modes
- 8-channel interrupt-driven single-cycle data transfer facilities via peripheral event controller (PEC)
- 16-priority-level interrupt system with 28 sources, sample-rate down to 50 ns
- Two multi-functional general purpose timer units with 5 timers
- Two serial channels (synchronous/asynchronous and high-speed-synchronous)
- Programmable watchdog timer
- Up to 77 general purpose I/O lines
- Supported by a wealth of development tools like C-compilers, macro-assembler packages, emulators, evaluation boards, HLL-debuggers, simulators, logic analyser disassemblers, programming boards
- On-chip bootstrap loader
- 100-Pin TQFP and PQFP package (EIAJ)

**PQFP100****TQFP100**

(Ordering Information at the end of the datasheet)

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## TABLE OF CONTENTS

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<b>ST10R165</b> .....	1
1 INTRODUCTION .....	3
2 FUNCTIONAL DESCRIPTION .....	10
3 MEMORY ORGANIZATION .....	11
4 EXTERNAL BUS CONTROLLER .....	12
5 CENTRAL PROCESSING UNIT (CPU) .....	13
6 INTERRUPT SYSTEM .....	15
7 GENERAL PURPOSE TIMER (GPT) UNIT .....	18
8 PARALLEL PORTS .....	21
9 SERIAL CHANNELS .....	21
10 WATCHDOG TIMER .....	22
11 INSTRUCTION SET SUMMARY .....	23
12 SPECIAL FUNCTION REGISTER OVERVIEW .....	25
13 ELECTRICAL CHARACTERISTICS .....	29
13.1 ABSOLUTE MAXIMUM RATINGS .....	29
13.2 PARAMETER INTERPRETATION .....	29
13.3 DC CHARACTERISTICS .....	30
13.4 TESTING WAVEFORMS .....	33
13.5 AC CHARACTERISTICS .....	34
14 PACKAGE MECHANICAL DATA .....	52
15 ORDERING INFORMATION .....	53

### 1 INTRODUCTION

The ST10R165 is a new derivative of the SGS-THOMSON ST10 family of full featured single-chip CMOS microcontrollers. It combines high CPU performance (up to 10 million instructions per second) with high peripheral functionality and enhanced IO-capabilities.

Figure 1. Logic Symbol

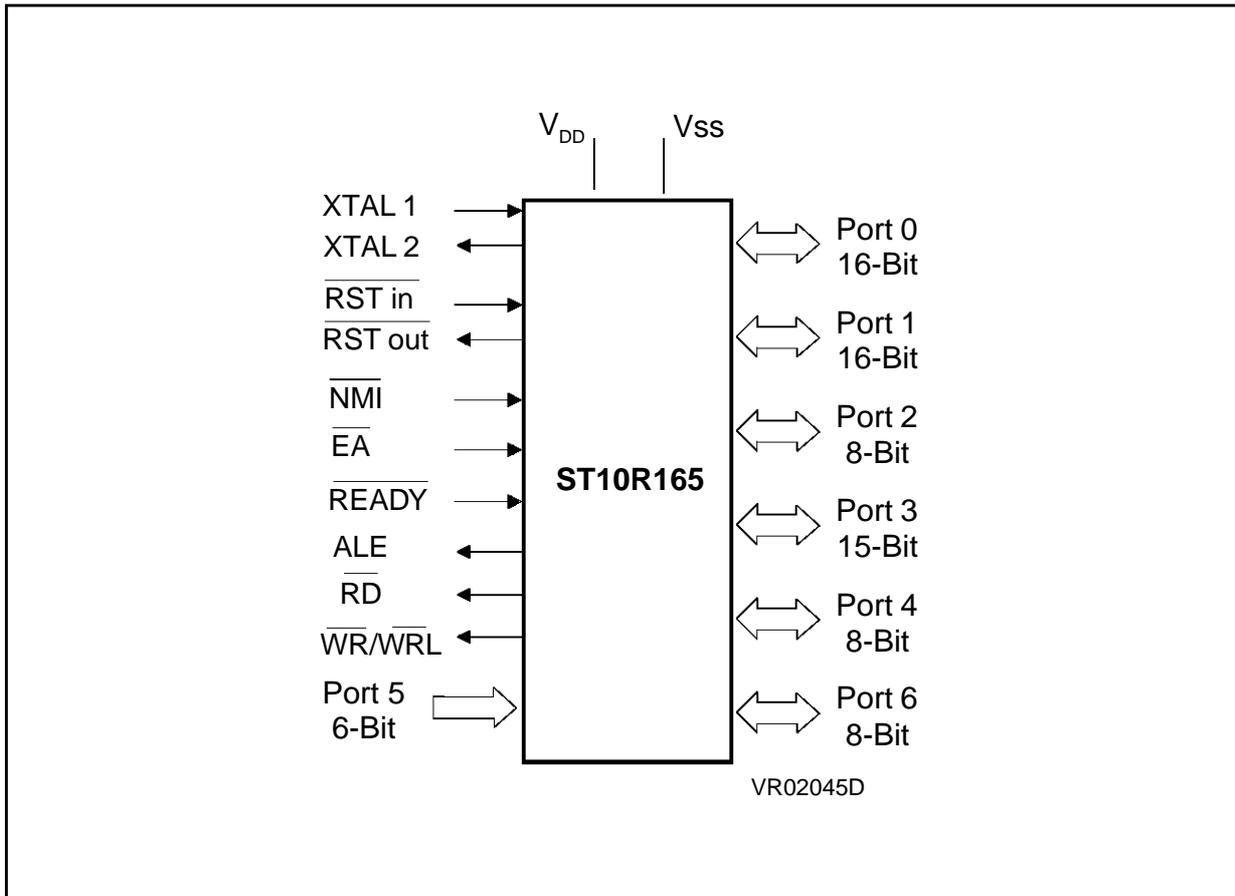


Figure 2. PQFP Pin Configuration (top view)

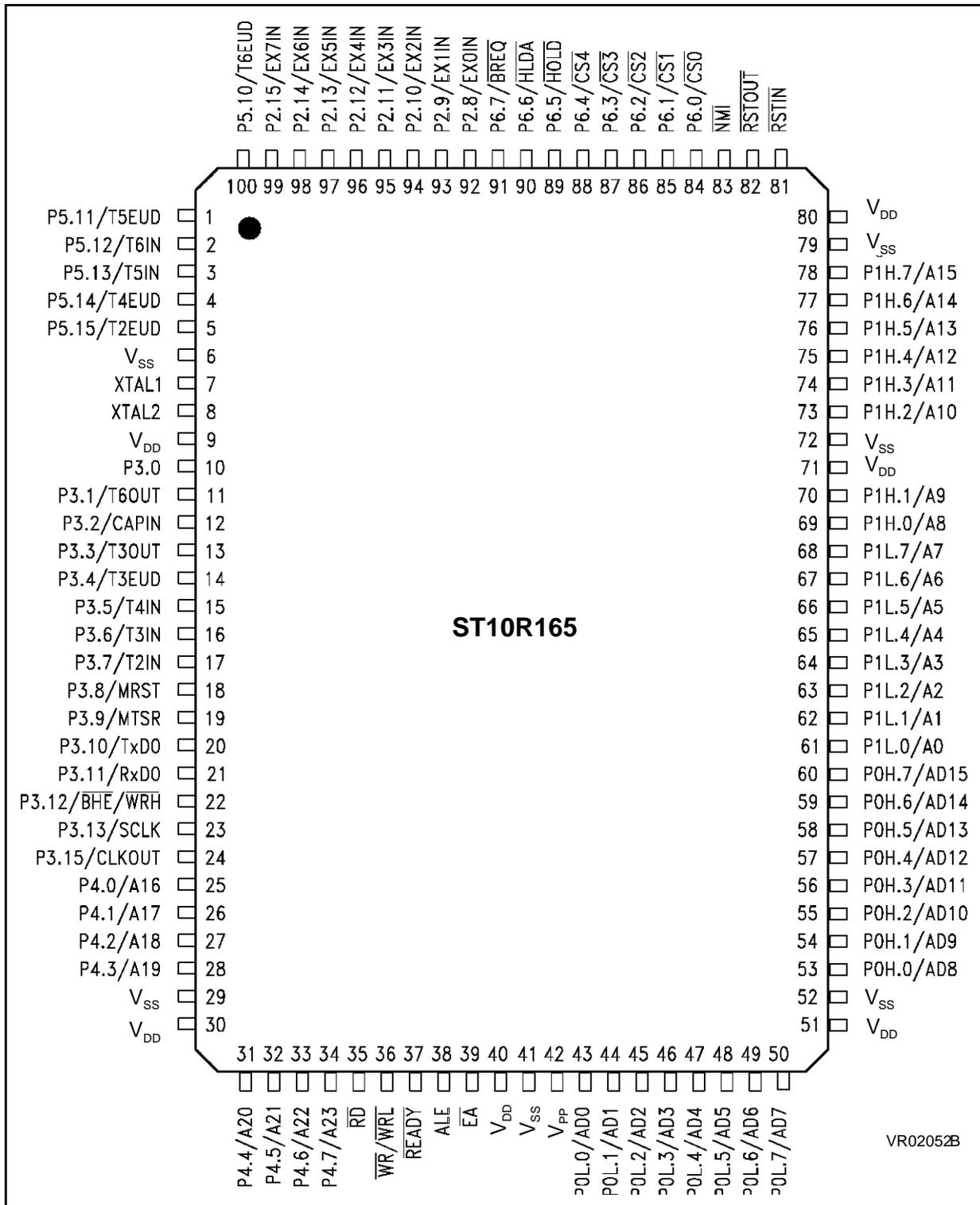




Table 1. Pin Definitions and Functions

Symbol	Pin Number		Input (I) Output (O)	Function
	PQFP	TQFP		
P5.10 – P5.15	100 1 - 5 100 1 2 3 4 5	98 99 - 3 98 99 100 1 2 3	I I I I I I I I	Port 5 is a 6-bit input-only port with Schmitt-Trigger characteristics. The pins of Port 5 also serve as timer inputs: P5.10 T6EUD GPT2 Timer T6 Ext.Up/Down Ctrl.Input P5.11 T5EUD GPT2 Timer T5 Ext.Up/Down Ctrl.Input P5.12 T6IN GPT2 Timer T6 Count Input P5.13 T5IN GPT2 Timer T5 Count Input P5.14 T4EUD GPT1 Timer T4 Ext.Up/Down Ctrl.Input P5.15 T2EUD GPT1 Timer T2 Ext.Up/Down Ctrl.Input
XTAL1  XTAL2	7  8	5  6	I  O	XTAL1: Input to the oscillator amplifier and input to the internal clock generator XTAL2: Output of the oscillator amplifier circuit. To clock the device from an external source, drive XTAL1, while leaving XTAL2 unconnected. Minimum and maximum high/low and rise/fall times specified in the AC Characteristics must be observed.
P3.0 – P3.13, P3.15	10 – 23, 24  11 12 13 14 15  16 17  18 19 20 21 22  23 24	8 21 22  9 10 11 12 13  14 15  16 17 18 19 20 21 22	I/O I/O I/O  O I O I I  I I  I/O I/O O I/O O O I/O O	Port 3 is a 15-bit (P3.14 is missing) bidirectional I/O port. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. Port 3 outputs can be configured as push/pull or open drain drivers. The following Port 3 pins also serve for alternate functions: P3.1 T6OUT GPT2 Timer T6 Toggle Latch Output P3.2 CAPIN GPT2 Register CAPREL Capture Input P3.3 T3OUT GPT1 Timer T3 Toggle Latch Output P3.4 T3EUD GPT1 Timer T3 Ext.Up/Down Ctrl.Input P3.5 T4IN GPT1 Timer T4 Input for Count/Gate/Reload/Capture P3.6 T3IN GPT1 Timer T3 Count/Gate Input P3.7 T2IN GPT1 Timer T2 Input for Count/Gate/Reload/Capture P3.8 MRST SSC Master-Rec./Slave-Transmit I/O P3.9 MTSR SSC Master-Transmit/Slave-Rec. O/I P3.10 TxD0 ASC0 Clock/Data Output (Asyn./Syn.) P3.11 RxD0 ASC0 Data Input (Asyn.) or I/O (Syn.) P3.12 BHE Ext. Memory High Byte Enable Signal, WRH Ext. Memory High Byte Write Strobe P3.13 SCLK SSC Master Clock Outp./Slave Cl. Inp. P3.15 CLKOUT System Clock Output (=CPU Clock)

Table 1. Pin Definitions and Functions (Cont'd)

Symbol	Pin Number		Input (I) Output (O)	Function
	PQFP	TQFP		
P4.0 – P4.7	25 - 28, 31 - 34	23 - 26, 29 - 32	I/O	Port 4 is an 8-bit bidirectional I/O port. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. In case of an external bus configuration, Port 4 can be used to output the segment address lines: P4.0 A16 Least Significant Segment Addr. Line ... .. P4.7 A23 Most Significant Segment Addr. Line
$\overline{RD}$	35	33	O	External Memory Read Strobe. $\overline{RD}$ is activated for every external instruction or data read access.
$\overline{WR}/\overline{WRL}$	36	34	O	External Memory Write Strobe. In $\overline{WR}$ -mode this pin is activated for every external data write access. In $\overline{WRL}$ -mode this pin is activated for low byte data write accesses on a 16-bit bus, and for every data write access on an 8-bit bus. See WRCFG in register SYSCON for mode selection.
READY	37	35	I	Ready Input. When the Ready function is enabled, a high level at this pin during an external memory access will force the insertion of memory cycle time waitstates until the pin returns to a low level.
ALE	38	36	O	Address Latch Enable Output. Can be used for latching the address into external memory or an address latch in the multiplexed bus modes.
EA	39	37	I	External Access Enable pin. A low level at this pin during and after Reset forces the ST10R165 to begin instruction execution out of external memory. A high level forces execution out of the internal ROM. The ST10R165 must have this pin tied to '0'.

Table 1. Pin Definitions and Functions (Cont'd)

Symbol	Pin Number		Input (I) Output (O)	Function																		
	PQFP	TQFP																				
PORT0: P0L.0 – P0L.7, P0H.0 - P0H.7	43 – 50 53 – 60	41 48 51 58	I/O	<p>PORT0 consists of the two 8-bit bidirectional I/O ports P0L and P0H. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state.</p> <p>In case of an external bus configuration, PORT0 serves as the address (A) and address/data (AD) bus in multiplexed bus modes and as the data (D) bus in demultiplexed bus modes.</p> <p><b>Demultiplexed bus modes:</b></p> <table> <tr> <td>Data Path Width:</td> <td>8-bit</td> <td>16-bit</td> </tr> <tr> <td>P0L.0 – P0L.7:</td> <td>D0 – D7</td> <td>D0 - D7</td> </tr> <tr> <td>P0H.0 – P0H.7:</td> <td>I/O</td> <td>D8 - D15</td> </tr> </table> <p><b>Multiplexed bus modes:</b></p> <table> <tr> <td>Data Path Width:</td> <td>8-bit</td> <td>16-bit</td> </tr> <tr> <td>P0L.0 – P0L.7:</td> <td>AD0 – AD7</td> <td>AD0 - AD7</td> </tr> <tr> <td>P0H.0 – P0H.7:</td> <td>A8 - A15</td> <td>AD8 - AD15</td> </tr> </table>	Data Path Width:	8-bit	16-bit	P0L.0 – P0L.7:	D0 – D7	D0 - D7	P0H.0 – P0H.7:	I/O	D8 - D15	Data Path Width:	8-bit	16-bit	P0L.0 – P0L.7:	AD0 – AD7	AD0 - AD7	P0H.0 – P0H.7:	A8 - A15	AD8 - AD15
Data Path Width:	8-bit	16-bit																				
P0L.0 – P0L.7:	D0 – D7	D0 - D7																				
P0H.0 – P0H.7:	I/O	D8 - D15																				
Data Path Width:	8-bit	16-bit																				
P0L.0 – P0L.7:	AD0 – AD7	AD0 - AD7																				
P0H.0 – P0H.7:	A8 - A15	AD8 - AD15																				
PORT1: P1L.0 – P1L.7, P1H.0 - P1H.7	61 - 68 69 - 70, 73 - 78	59 66 67 - 68 71 - 76	I/O	<p>PORT1 consists of the two 8-bit bidirectional I/O ports P1L and P1H. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. PORT1 is used as the 16-bit address bus (A) in demultiplexed bus modes and also after switching from a demultiplexed bus mode to a multiplexed bus mode.</p>																		
RSTIN	81	79	I	<p>Reset Input with Schmitt-Trigger characteristics. A low level at this pin for a specified duration while the oscillator is running resets the ST10R165. An internal pullup resistor permits power-on reset using only a capacitor connected to <math>V_{SS}</math>.</p>																		
RSTOUT	82	80	O	<p>Internal Reset Indication Output. This pin is set to a low level when the part is executing either a hardware-, a software- or a watchdog timer reset. RSTOUT remains low until the EINIT (end of initialization) instruction is executed.</p>																		
$\overline{NMI}$	83	81	I	<p>Non-Maskable Interrupt Input. A high to low transition at this pin causes the CPU to vector to the NMI trap routine. When the PWRDN (power down) instruction is executed, the <math>\overline{NMI}</math> pin must be low in order to force the ST10R65 to go into power down mode. If <math>\overline{NMI}</math> is high, when PWRDN is executed, the part will continue to run in normal mode.</p> <p>If not used, pin <math>\overline{NMI}</math> should be pulled high externally.</p>																		

Table 1. Pin Definitions and Functions (Cont'd)

Symbol	Pin Number		Input (I) Output (O)	Function
	PQFP	TQFP		
P6.0 – P6.7	84 - 91	82 - 89	I/O	<p>Port 6 is an 8-bit bidirectional I/O port. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. Port 6 outputs can be configured as push/pull or open drain drivers.</p> <p>The following Port 6 pins also serve for alternate functions:</p> <p>P6.0 <math>\overline{CS0}</math> Chip Select 0 Output</p> <p>...     ...     ...</p> <p>P6.4 <math>\overline{CS4}</math> Chip Select 4 Output</p> <p>P6.5 <math>\overline{HOLD}</math> External Master Hold Request Input</p> <p>P6.6 <math>\overline{HLDA}</math> Hold Acknowledge Output</p> <p>P6.7 <math>\overline{BREQ}</math> Bus Request Output</p>
P2.8 – P2.15	92 - 99	90 - 97	I/O	<p>Port 2 is an 8-bit bidirectional I/O port. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. Port 2 outputs can be configured as push/pull or open drain drivers.</p> <p>The following Port 2 pins also serve for alternate functions:</p> <p>P2.8 EX0IN Fast External Interrupt 0 Input</p> <p>...     ...     ...</p> <p>P2.15 EX7IN Fast External Interrupt 7 Input</p>
V <sub>PP</sub>	42	40	-	<p>Flash programming voltage. This pin accepts the programming voltage for flash versions of the ST10R165.</p> <p><b>Note:</b> This pin is not connected (<b>NC</b>) on non-flash versions.</p>
V <sub>DD</sub>	9, 30, 40, 51, 71, 80	7, 28, 38, 49, 69, 78	-	<p>Digital Supply Voltage: + 5 V during normal operation and idle mode. ≥ 2.5 V during power down mode</p>
V <sub>SS</sub>	6, 29, 41, 52, 72, 79	4, 27, 39, 50, 70, 77	-	Digital Ground.



### **3 MEMORY ORGANIZATION**

The memory space of the ST10R165 is configured in a Von Neumann architecture which means that code memory, data memory, registers and I/O ports are organized within the same linear address space which includes 16 MBytes. The entire memory space can be accessed bitwise or wordwise. Particular portions of the on-chip memory have additionally been made directly bit addressable. The ST10R165 is prepared to incorporate on-chip mask-programmable ROM for code or constant data. Currently no ROM is integrated.

2 KBytes of on-chip RAM are provided as a storage for user defined variables, for the system stack, general purpose register banks and even for code. A register bank can consist of up to 16 wordwide (R0 to R15) and/or bitwise (RL0, RH0, ..., RL7, RH7) named General Purpose Registers (GPRs).

1024 bytes (2 \* 512 bytes) of the address space are reserved for the Special Function Register areas (SFR space and ESFR space). SFRs are wordwide registers which are used for controlling and monitoring functions of the different on-chip units. Unused SFR addresses are reserved for future derivatives of the ST10R165.

In order to meet the needs of designs where more memory is required than is provided on chip, up to 16 MBytes of external RAM and/or ROM can be connected to the microcontroller.

## 4 EXTERNAL BUS CONTROLLER

All of the external memory accesses are performed by a particular on-chip External Bus Controller (EBC). It can be programmed either to Single Chip Mode when no external memory is required, or to one of four different external memory access modes, which are as follows:

- 16-/18-/20-/24-bit Addresses, 16-bit Data, Demultiplexed
- 16-/18-/20-/24-bit Addresses, 16-bit Data, Multiplexed
- 16-/18-/20-/24-bit Addresses, 8-bit Data, Multiplexed
- 16-/18-/20-/24-bit Addresses, 8-bit Data, Demultiplexed

In the demultiplexed bus modes, addresses are output on PORT1 and data is input/output on PORT0. In the multiplexed bus modes both addresses and data use PORT0 for input/output.

Important timing characteristics of the external bus interface (Memory Cycle Time, Memory Tri-State Time, Length of ALE and Read Write Delay) have been made programmable to allow the user the adoption of a wide range of different types of memories. In addition, different address ranges may be accessed with different bus characteristics. Up to 5 external  $\overline{CS}$  signals can be generated in order to save external glue logic. Access to very slow memories is supported via a particular 'Ready' function. A  $\overline{HOLD}/\overline{HLDA}$  protocol is available for bus arbitration.

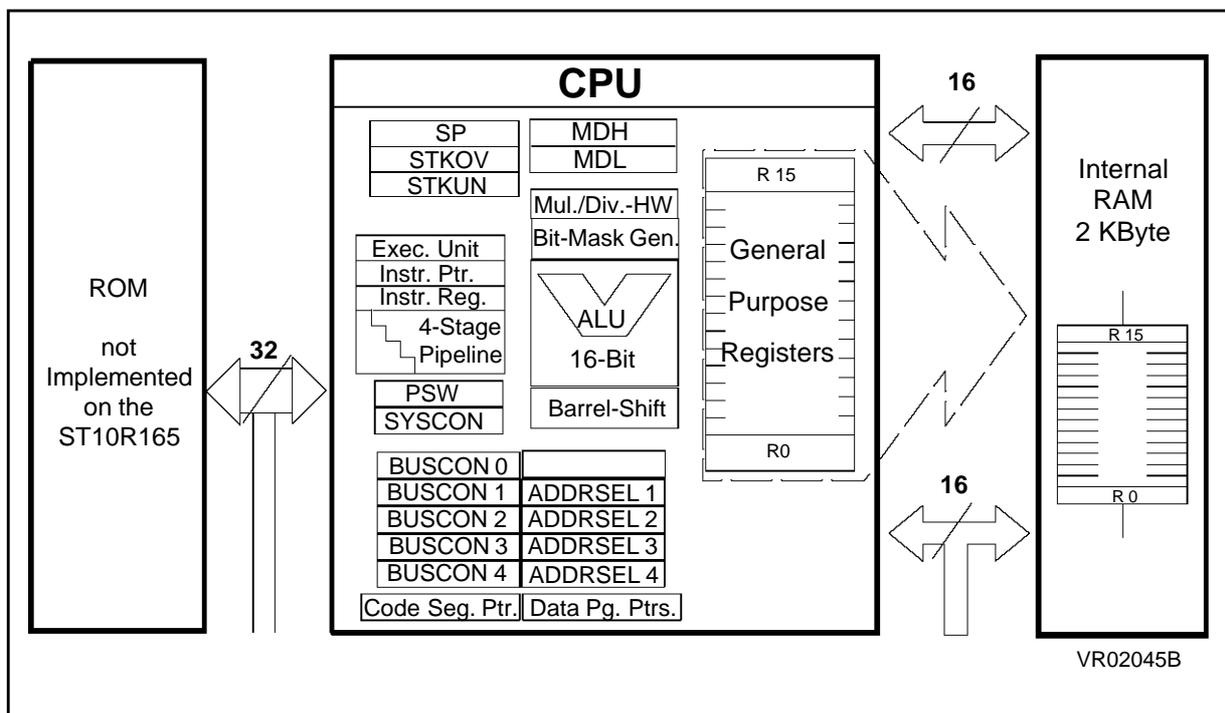
For applications which require less than 16 MBytes of external memory space, this address space can be restricted to 1 MByte, 256 KByte or to 64 KByte. In this case Port 4 outputs four, two or no address lines at all. It outputs all 8 address lines, if an address space of 16 MBytes is used.

### 5 CENTRAL PROCESSING UNIT (CPU)

The main core of the CPU consists of a 4-stage instruction pipeline, a 16-bit arithmetic and logic unit (ALU) and dedicated SFRs. Additional hardware has been added for a separate multiply and divide unit, a bit-mask generator and a barrel shifter.

Based on these hardware provisions, most of the ST10R165's instructions can be executed in just one machine cycle which requires 100 ns at 20 MHz CPU clock. For example, shift and rotate instructions are always processed during one machine cycle independent of the number of bits to be shifted. All multiple-cycle instructions have been optimized so that they can be executed very fast as well: branches in 2 cycles, a 16 × 16 bit multiplication in 5 cycles and a 32/16 bit division in 10 cycles. Another pipeline optimization, the named 'Jump Cache', allows reducing the execution time of repeatedly performed jumps in a loop from 2 cycles to 1 cycle.

Figure 5. CPU Block Diagram



### Central Processing Unit (Cont'd)

The CPU includes an actual register context consisting of up to 16 worldwide GPRs which are physically allocated within the on-chip RAM area. A Context Pointer (CP) register determines the base address of the active register bank to be accessed by the CPU at a time. The number of register banks is only restricted by the available internal RAM space. For easy parameter passing, a register bank may overlap others.

A system stack of up to 2048 bytes is provided as a storage for temporary data. The system stack is allocated in the on-chip RAM area, and it is accessed by the CPU via the stack pointer (SP) register. Two separate SFRs, STKOV and STKUN, are implicitly compared against the stack pointer value upon each stack access for the detection of a stack overflow or underflow.

The high performance offered by the hardware implementation of the CPU can efficiently be utilized by a programmer via the highly efficient ST10R165 instruction set which includes the following instruction classes:

- Arithmetic Instructions
- Logical Instructions
- Boolean Bit Manipulation Instructions
- Compare and Loop Control Instructions
- Shift and Rotate Instructions
- Prioritize Instruction
- Data Movement Instructions
- System Stack Instructions
- Jump and Call Instructions
- Return Instructions
- System Control Instructions
- Miscellaneous Instructions

The basic instruction length is either 2 or 4 bytes. Possible operand types are bits, bytes and words. A variety of direct, indirect or immediate addressing modes are provided to specify the required operands.

## 6 INTERRUPT SYSTEM

With an interrupt response time within a range from just 250 ns to 600 ns (in case of internal program execution), the ST10R165 is capable of reacting very fast to the occurrence of non-deterministic events.

The architecture of the ST10R165 supports several mechanisms for fast and flexible response to service requests that can be generated from various sources internal or external to the microcontroller. Any of these interrupt requests can be programmed to being serviced by the Interrupt Controller or by the Peripheral Event Controller (PEC).

In contrast to a standard interrupt service where the current program execution is suspended and a branch to the interrupt vector table is performed, just one cycle is 'stolen' from the current CPU activity to perform a PEC service. A PEC service implies a single byte or word data transfer between any two memory locations with an additional increment of either the PEC source or the destination pointer. An individual PEC transfer counter is implicitly decremented for each PEC service except when performing in the continuous transfer mode. When this counter reaches zero, a standard interrupt is performed to the corresponding source related vector location. PEC services are very well suited, for example, for supporting the transmission or reception of blocks of data. The ST10R165 has 8 PEC channels each of which offers such fast interrupt-driven data transfer capabilities.

A separate control register which contains an interrupt request flag, an interrupt enable flag and an interrupt priority bitfield exists for each of the possible interrupt sources. Via its related register, each source can be programmed to one of sixteen interrupt priority levels. Once having been accepted by the CPU, an interrupt service can only be interrupted by a higher prioritized service request. For the standard interrupt processing, each of the possible interrupt sources has a dedicated vector location.

Fast external interrupt inputs are provided to service external interrupts with high precision requirements. These fast interrupt inputs feature programmable edge detection (rising edge, falling edge or both edges).

Software interrupts are supported by means of the 'TRAP' instruction in combination with an individual trap (interrupt) number.

The following table shows all of the possible ST10R165 interrupt sources and the corresponding hardware-related interrupt flags, vectors, vector locations and trap (interrupt) numbers:

**Note:** Four nodes in the table (X-Peripheral nodes) are prepared to accept interrupt requests from integrated X-Bus peripherals. Nodes, where no X-Peripherals are connected, may be used to generate software controlled interrupt requests by setting the respective XPnIR bit. Also the three listed Software Nodes can be used for this purpose.

## Interrupt System (Cont'd)

Source of Interrupt or PEC Service Request	Request Flag	Enable Flag	Interrupt Vector	Vector Location	Trap Number
External Interrupt 0	CC8IR	CC8IE	CC8INT	60h	18h
External Interrupt 1	CC9IR	CC9IE	CC9INT	64h	19h
External Interrupt 2	CC10IR	CC10IE	CC10INT	68h	1Ah
External Interrupt 3	CC11IR	CC11IE	CC11INT	6Ch	1Bh
External Interrupt 4	CC12IR	CC12IE	CC12INT	70h	1Ch
External Interrupt 5	CC13IR	CC13IE	CC13INT	74h	1Dh
External Interrupt 6	CC14IR	CC14IE	CC14INT	78h	1Eh
External Interrupt 7	CC15IR	CC15IE	CC15INT	7Ch	1Fh
GPT1 Timer 2	T2IR	T2IE	T2INT	88h	22h
GPT1 Timer 3	T3IR	T3IE	T3INT	8Ch	23h
GPT1 Timer 4	T4IR	T4IE	T4INT	90h	24h
GPT2 Timer 5	T5IR	T5IE	T5INT	94h	25h
GPT2 Timer 6	T6IR	T6IE	T6INT	98h	26h
GPT2 CAPREL Register	CRIR	CRIE	CRINT	9Ch	27h
ASC0 Transmit	S0TIR	S0TIE	S0TINT	A8h	2Ah
ASC0 Transmit Buffer	S0TBIR	S0TBIE	S0TBINT	11Ch	47h
ASC0 Receive	S0RIR	S0RIE	S0RINT	ACh	2Bh
ASC0 Error	S0EIR	S0EIE	S0EINT	B0h	2Ch
SSC Transmit	SCTIR	SCTIE	SCTINT	B4h	2Dh
SSC Receive	SCRIR	SCRIE	SCRINT	B8h	2Eh
SSC Error	SCEIR	SCEIE	SCEINT	BCh	2Fh
X-Peripheral Node 0	XP0IR	XP0IE	XP0INT	100h	40h
X-Peripheral Node 1	XP1IR	XP1IE	XP1INT	104h	41h
X-Peripheral Node 2	XP2IR	XP2IE	XP2INT	108h	42h
X-Peripheral Node 3	XP3IR	XP3IE	XP3INT	10Ch	43h
Software Node	CC29IR	CC29IE	CC29INT	110h	44h
Software Node	CC30IR	CC30IE	CC30INT	114h	45h
Software Node	CC31IR	CC31IE	CC31INT	118h	46h

### Interrupt System (Cont'd)

The ST10R165 also provides an excellent mechanism to identify and to process exceptions or error conditions that arise during run-time, named 'Hardware Traps'. Hardware traps cause immediate non-maskable system reaction which is similar to a standard interrupt service (branching to a dedicated vector table location). The occurrence of a hardware trap is additionally signified by an individual bit in the trap flag register (TFR). Except when another higher prioritized trap service is in progress, a hardware trap will interrupt any actual program execution. In turn, hardware trap services can normally not be interrupted by standard or PEC interrupts.

The following table shows all of the possible exceptions or error conditions that can arise during run-time:

Exception Condition	Trap Flag	Trap Vector	Vector Location	Trap Number	Trap Priority
Reset Functions: Hardware Reset Software Reset Watchdog Timer Overflow		RESET RESET RESET	00'0000h 00'0000h 00'0000h	00h 00h 00h	III III III
Class A Hardware Traps: Non-Maskable Interrupt Stack Overflow Stack Underflow	NMI STKOF STKUF	NMITRAP STOTRAP STUTRAP	00'0008h 00'0010h 00'0018h	02h 04h 06h	II II II
Class B Hardware Traps: Undefined Opcode Protected Instruction Fault Illegal Word Operand Access Illegal Instruction Access Illegal External Bus Access	UNDOPC PRTFLT  ILLOPA  ILLINA ILLBUS	BTRAP BTRAP  BTRAP  BTRAP BTRAP	00'0028h 00'0028h  00'0028h  00'0028h 00'0028h	0Ah 0Ah  0Ah  0Ah 0Ah	I I  I  I I
Reserved			[2Ch – 3Ch]	[0Bh – 0Fh]	
Software Traps TRAP Instruction			Any [00'0000h – 00'01FCh] in steps of 4h	Any [00h – 7Fh]	Current CPU Prior- ity

## 7 GENERAL PURPOSE TIMER (GPT) UNIT

The GPT unit represents a very flexible multifunctional timer/counter structure which may be used for many different time related tasks such as event timing and counting, pulse width and duty cycle measurements, pulse generation, or pulse multiplication.

The GPT unit incorporates five 16-bit timers which are organized in two separate modules, GPT1 and GPT2. Each timer in each module may operate independently in a number of different modes, or may be concatenated with another timer of the same module.

Each of the three timers T2, T3, T4 of module GPT1 can be configured individually for one of three basic modes of operation, which are Timer, Gated Timer, and Counter Mode. In Timer Mode, the input clock for a timer is derived from the CPU clock, divided by a programmable prescaler, while Counter Mode allows a timer to be clocked in reference to external events.

Pulse width or duty cycle measurement is supported in Gated Timer Mode, where the operation of a timer is controlled by the 'gate' level on an external input pin. For these purposes, each timer has one associated port pin (TxIN) which serves as gate or clock input. The maximum resolution of the timers in module GPT1 is 400 ns (@ 20 MHz CPU clock).

The count direction (up/down) for each timer is programmable by software or may additionally be altered dynamically by an external signal on a port pin (TxEUD) to facilitate e. g. position tracking.

Timers T3 and T4 have output toggle latches (TxOTL) which change their state on each timer over-flow/underflow. The state of these latches may be output on port pins (TxOUT) e. g. for time out monitoring of external hardware components, or may be used internally to clock timers T2 and T4 for measuring long time periods with high resolution.

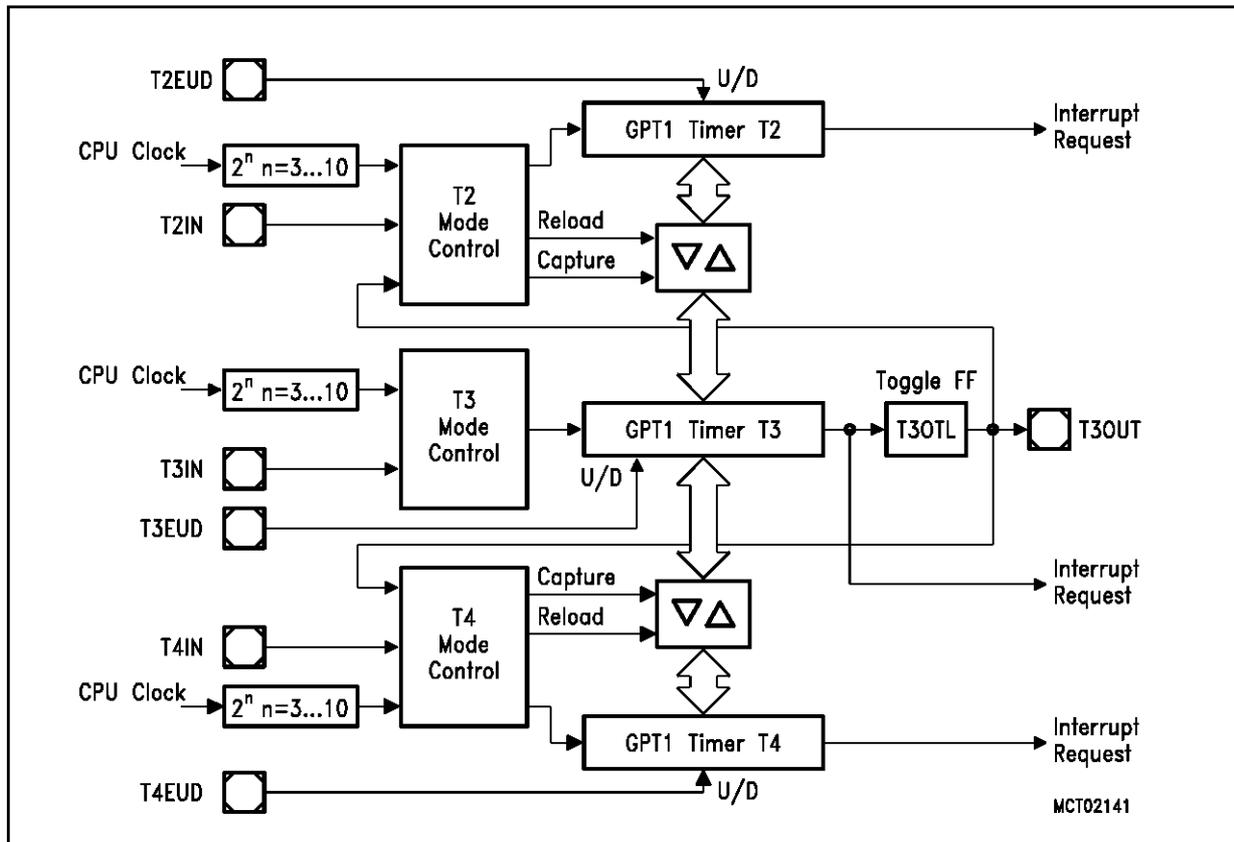
In addition to their basic operating modes, timers T2 and T4 may be configured as reload or capture registers for timer T3. When used as capture or reload registers, timers T2 and T4 are stopped. The contents of timer T3 is captured into T2 or T4 in response to a signal at their associated input pins (TxIN). Timer T3 is reloaded with the contents of T2 or T4 triggered either by an external signal or by a selectable state transition of its toggle latch T3OTL. When both T2 and T4 are configured to alternately reload T3 on opposite state transitions of T3OTL with the low and high times of a PWM signal, this signal can be constantly generated without software intervention.

**General Purpose Timer (GPT) Unit (Cont'd)**

With its maximum resolution of 200 ns (@ 20 MHz), the GPT2 module provides precise event control and time measurement. It includes two timers (T5, T6) and a capture/reload register (CAPREL). Both timers can be clocked with an input clock which is derived from the CPU clock via a programmable prescaler or with external signals. The count direction (up/down) for each timer is programmable by software or may additionally be altered dynamically by an external signal on a port pin (TxEUD). Concatenation of the timers is supported via the output toggle latch (T6OTL) of timer T6, which changes its state on each timer overflow/underflow.

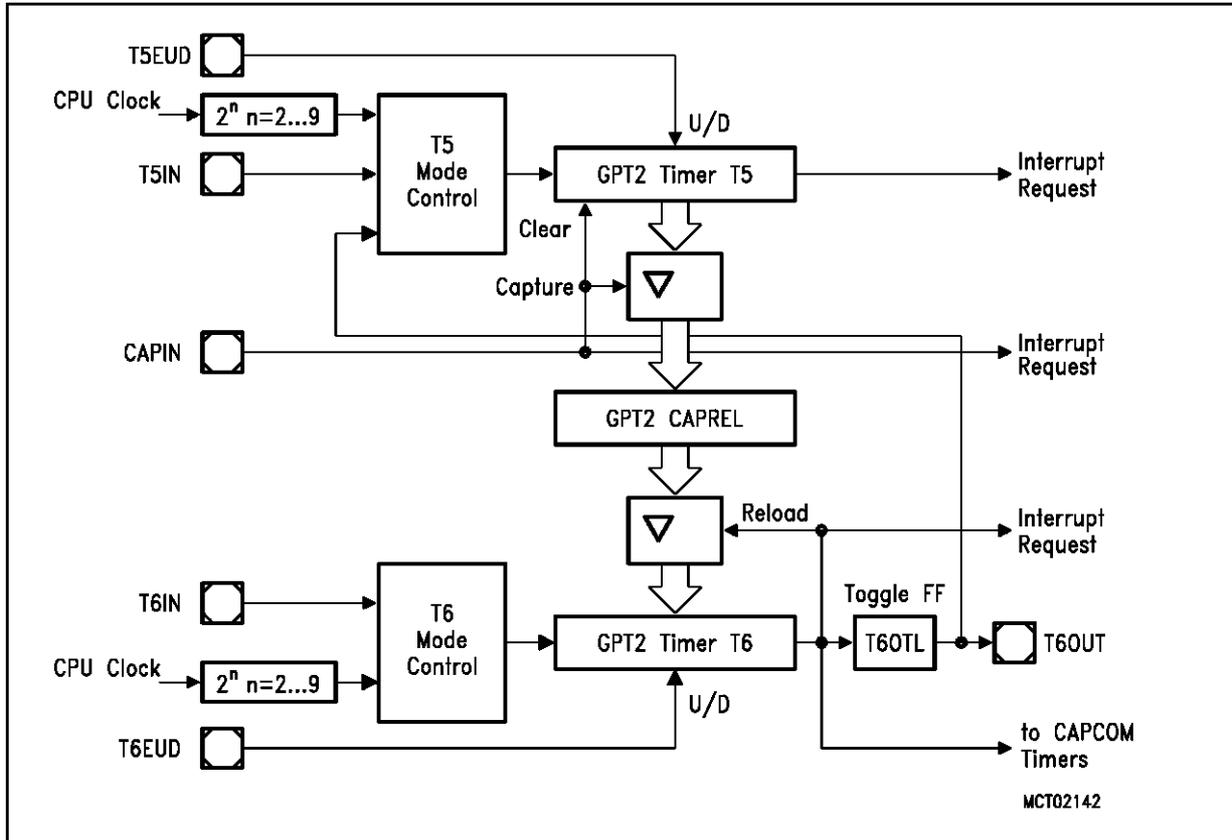
The state of this latch may be used to clock timer T5, or it may be output on a port pin (T6OUT). The overflows/underflows of timer T6 can additionally be used to clock the CAPCOM timers T0 or T1, and to cause a reload from the CAPREL register. The CAPREL register may capture the contents of timer T5 based on an external signal transition on the corresponding port pin (CAPIN), and timer T5 may optionally be cleared after the capture procedure. This allows absolute time differences to be measured or pulse multiplication to be performed without software overhead.

**Figure 6. Block Diagram of GPT1**



General Purpose Timer (GPT) Unit (Cont'd)

Figure 7. Block Diagram of GPT2



## **8 PARALLEL PORTS**

The ST10R165 provides up to 77 I/O lines which are organized into six input/output ports and one input port. All port lines are bit-addressable, and all input/output lines are individually (bit-wise) programmable as inputs or outputs via direction registers. The I/O ports are true bidirectional ports which are switched to high impedance state when configured as inputs. The output drivers of three I/O ports can be configured (pin by pin) for push/pull operation or open-drain operation via control registers. During the internal reset, all port pins are configured as inputs.

All port lines have programmable alternate input or output functions associated with them. PORT0 and PORT1 may be used as address and data lines when accessing external memory, while Port 4 outputs the additional segment address bits A23/19/17...A16 in systems where segmentation is enabled to access more than 64KBytes of memory. Port 6 provides optional bus arbitration signals ( $\overline{BREQ}$ ,  $\overline{HLDA}$ ,  $\overline{HOLD}$ ) and chip select signals. Port 3 includes alternate functions of timers, serial interfaces, the optional bus control signal  $\overline{BHE}$  and the system clock output (CLKOUT). Port 5 is used for timer control signals. All port lines that are not used for these alternate functions may be used as general purpose I/O lines.

## **9 SERIAL CHANNELS**

Serial communication with other microcontrollers, processors, terminals or external peripheral components is provided by two serial interfaces with different functionality, an Asynchronous/Synchronous Serial Channel (ASC0) and a High-Speed Synchronous Serial Channel (SSC).

They support full-duplex asynchronous communication up to 625 Kbaud and half-duplex synchronous communication up to 5 Mbaud (2.5 Mbaud on the ASC0) @ 20 MHz system clock.

Two dedicated baud rate generators allow to set up all standard baud rates without oscillator tuning. For transmission, reception, and erroneous reception 3 separate interrupt vectors are provided for each serial channel.

In asynchronous mode, 8- or 9-bit data frames are transmitted or received, preceded by a start bit and terminated by one or two stop bits. For multiprocessor communication, a mechanism to distinguish address from data bytes has been included (8-bit data + wake up bit mode).

### Serial Channels (Cont'd)

In synchronous mode, the ASC0 transmits or receives bytes (8 bits) synchronously to a shift clock which is generated by the ASC0. The SSC transmits or receives characters of 2...16 bits length synchronously to a shift clock which can be generated by the SSC (master mode) or by an external master (slave mode). The SSC can start shifting with the LSB or with the MSB, while the ASC0 always shifts the LSB first.

A loop back option is available for testing purposes.

A number of optional hardware error detection capabilities has been included to increase the reliability of data transfers. A parity bit can automatically be generated on transmission or be checked on reception. Framing error detection allows to recognize data frames with missing stop bits. An overrun error will be generated, if the last character received has not been read out of the receive buffer register at the time the reception of a new character is complete.

## 10 WATCHDOG TIMER

The Watchdog Timer represents one of the fail-safe mechanisms which have been implemented to prevent the controller from malfunctioning for longer periods of time.

The Watchdog Timer is always enabled after a reset of the chip, and can only be disabled in the time interval until the EINIT (end of initialization) instruction has been executed. Thus, the chip's start-up procedure is always monitored. The software has to be designed to service the Watchdog Timer before it overflows. If, due to hardware or software related failures, the software fails to do so, the Watchdog Timer overflows and generates an internal hardware reset and pulls the  $\overline{\text{RSTOUT}}$  pin low in order to allow external hardware components to be reset.

The Watchdog Timer is a 16-bit timer, clocked with the system clock divided either by 2 or by 128. The high byte of the Watchdog Timer register can be set to a pre-specified reload value (stored in WDTREL) in order to allow further variation of the monitored time interval. Each time it is serviced by the application software, the high byte of the Watchdog Timer is reloaded. Thus, time intervals between 25  $\mu\text{s}$  and 420 ms can be monitored (@ 20 MHz). The default Watchdog Timer interval after reset is 6.55 ms (@ 20 MHz).

## 11 INSTRUCTION SET SUMMARY

The table below lists the instructions of the ST10R165 in a condensed way.

The various addressing modes that can be used with a specific instruction, the operation of the instructions, parameters for conditional execution of instructions, and the opcodes for each instruction can be found in the “**ST10 Programming Manual**”.

This document also provides a detailed description of each instruction.

Mnemonic	Description	Bytes
ADD(B)	Add word (byte) operands	2 / 4
ADDC(B)	Add word (byte) operands with Carry	2 / 4
SUB(B)	Subtract word (byte) operands	2 / 4
SUBC(B)	Subtract word (byte) operands with Carry	2 / 4
MUL(U)	(Un)Signed multiply direct GPR by direct GPR (16-16-bit)	2
DIV(U)	(Un)Signed divide register MDL by direct GPR (16-/16-bit)	2
DIVL(U)	(Un)Signed long divide reg. MD by direct GPR (32-/16-bit)	2
CPL(B)	Complement direct word (byte) GPR	2
NEG(B)	Negate direct word (byte) GPR	2
AND(B)	Bitwise AND, (word/byte operands)	2 / 4
OR(B)	Bitwise OR, (word/byte operands)	2 / 4
XOR(B)	Bitwise XOR, (word/byte operands)	2 / 4
BCLR	Clear direct bit	2
BSET	Set direct bit	2
BMOV(N)	Move (negated) direct bit to direct bit	4
BAND, BOR, BXOR	AND/OR/XOR direct bit with direct bit	4
BCMP	Compare direct bit to direct bit	4
BFLDH/L	Bitwise modify masked high/low byte of bit-addressable direct word memory with immediate data	4
CMP(B)	Compare word (byte) operands	2 / 4
CMPD1/2	Compare word data to GPR and decrement GPR by 1/2	2 / 4
CMPI1/2	Compare word data to GPR and increment GPR by 1/2	2 / 4
PRIOR	Determine number of shift cycles to normalize direct word GPR and store result in direct word GPR	2
SHL / SHR	Shift left/right direct word GPR	2
ROL / ROR	Rotate left/right direct word GPR	2

## Instruction Set (Cont'd)

Mnemonic	Description	Bytes
ASHR	Arithmetic (sign bit) shift right direct word GPR	2
MOV(B)	Move word (byte) data	2 / 4
MOVBS	Move byte operand to word operand with sign extension	2 / 4
MOVBZ	Move byte operand to word operand. with zero extension	2 / 4
JMPA, JMPI, JMPR	Jump absolute/indirect/relative if condition is met	4
JMPS	Jump absolute to a code segment	4
J(N)B	Jump relative if direct bit is (not) set	4
JBC	Jump relative and clear bit if direct bit is set	4
JNBS	Jump relative and set bit if direct bit is not set	4
CALLA, CALLI, CALLR	Call absolute/indirect/relative subroutine if condition is met	4
CALLS	Call absolute subroutine in any code segment	4
PCALL	Push direct word register onto system stack and call absolute subroutine	4
TRAP	Call interrupt service routine via immediate trap number	2
PUSH, POP	Push/pop direct word register onto/from system stack	2
SCXT	Push direct word register onto system stack and update register with word operand	4
RET	Return from intra-segment subroutine	2
RETS	Return from inter-segment subroutine	2
RETP	Return from intra-segment subroutine and pop direct word register from system stack	2
RETI	Return from interrupt service subroutine	2
SRST	Software Reset	4
IDLE	Enter Idle Mode	4
PWRDN	Enter Power Down Mode (supposes NMI-pin being low)	4
SRVWDT	Service Watchdog Timer	4
DISWDT	Disable Watchdog Timer	4
EINIT	Signify End-of-Initialization on RSTOUT-pin	4
ATOMIC	Begin ATOMIC sequence	2
EXTR	Begin EXTENDED Register sequence	2
EXTP(R)	Begin EXTENDED Page (and Register) sequence	2 / 4
EXTS(R)	Begin EXTENDED Segment (and Register) sequence	2 / 4
NOP	Null operation	2

## 12 SPECIAL FUNCTION REGISTER OVERVIEW

The following table lists all SFRs which are implemented in the ST10R165 in alphabetical order.

**Bit-addressable** SFRs are marked with the letter “b” in column “Name”. SFRs within the **Extended SFR-Space** (ESFRs) are marked with the letter “E” in column “Physical Address”.

An SFR can be specified via its individual mnemonic name. Depending on the selected addressing mode, an SFR can be accessed via its physical address (using the Data Page Pointers), or via its short 8-bit address (without using the Data Page Pointers).

Name	Physical Address	8-Bit Address	Description	Reset Value
ADDRSEL1	FE18h	0Ch	Address Select Register 1	0000h
ADDRSEL2	FE1Ah	0Dh	Address Select Register 2	0000h
ADDRSEL3	FE1Ch	0Eh	Address Select Register 3	0000h
ADDRSEL4	FE1Eh	0Fh	Address Select Register 4	0000h
BUSCON0b	FF0Ch	86h	Bus Configuration Register 0	0XX0h
BUSCON1b	FF14h	8Ah	Bus Configuration Register 1	0000h
BUSCON2b	FF16h	8Bh	Bus Configuration Register 2	0000h
BUSCON3b	FF18h	8Ch	Bus Configuration Register 3	0000h
BUSCON4b	FF1Ah	8Dh	Bus Configuration Register 4	0000h
CAPREL	FE4Ah	25h	GPT2 Capture/Reload Register	0000h
CC8IC b	FF88h	C4h	EX0IN Interrupt Control Register	0000h
CC9IC b	FF8Ah	C5h	EX1IN Interrupt Control Register	0000h
CC10IC b	FF8Ch	C6h	EX2IN Interrupt Control Register	0000h
CC11IC b	FF8Eh	C7h	EX3IN Interrupt Control Register	0000h
CC12IC b	FF90h	C8h	EX4IN Interrupt Control Register	0000h
CC13IC b	FF92h	C9h	EX5IN Interrupt Control Register	0000h
CC14IC b	FF94h	CAh	EX6IN Interrupt Control Register	0000h
CC15IC b	FF96h	CBh	EX7IN Interrupt Control Register	0000h
CC29IC b	F184h <sub>E</sub>	C2h	Software Node Interrupt Control Register	0000h
CC30IC b	F18Ch <sub>E</sub>	C6h	Software Node Interrupt Control Register	0000h
CC31IC b	F194h <sub>E</sub>	CAh	Software Node Interrupt Control Register	0000h
CP	FE10h	08h	CPU Context Pointer Register	FC00h

### Special Function Register Overview (Cont'd)

Name	Physical Address	8-Bit Address	Description	Reset Value
CRIC	b FF6Ah	B5h	GPT2 CAPREL Interrupt Control Register	0000h
CSP	FE08h	04h	CPU Code Segment Pointer Register (read only)	0000h
DP0L	b F100h <sub>E</sub>	80h	P0L Direction Control Register	00h
DP0H	b F102h <sub>E</sub>	81h	P0h Direction Control Register	00h
DP1L	b F104h <sub>E</sub>	82h	P1L Direction Control Register	00h
DP1H	b F106h <sub>E</sub>	83h	P1h Direction Control Register	00h
DP2	b FFC2h	E1h	Port 2 Direction Control Register	0000h
DP3	b FFC6h	E3h	Port 3 Direction Control Register	0000h
DP4	b FFCAh	E5h	Port 4 Direction Control Register	00h
DP6	b FFCEh	E7h	Port 6 Direction Control Register	00h
DPP0	FE00h	00h	CPU Data Page Pointer 0 Register (10 bits)	0000h
DPP1	FE02h	01h	CPU Data Page Pointer 1 Register (10 bits)	0001h
DPP2	FE04h	02h	CPU Data Page Pointer 2 Register (10 bits)	0002h
DPP3	FE06h	03h	CPU Data Page Pointer 3 Register (10 bits)	0003h
EXICON	b F1C0h <sub>E</sub>	E0h	External Interrupt Control Register	0000h
MDC	b FF0Eh	87h	CPU Multiply Divide Control Register	0000h
MDH	FE0Ch	06h	CPU Multiply Divide Register – High Word	0000h
MDL	FE0Eh	07h	CPU Multiply Divide Register – Low Word	0000h
ODP2	b F1C2h <sub>E</sub>	E1h	Port 2 Open Drain Control Register	0000h
ODP3	b F1C6h <sub>E</sub>	E3h	Port 3 Open Drain Control Register	0000h
ODP6	b F1CEh <sub>E</sub>	E7h	Port 6 Open Drain Control Register	00h
ONES	FF1Eh	8Fh	Constant Value 1's Register (read only)	FFFFh
P0L	b FF00h	80h	Port 0 Low Register (Lower half of PORT0)	00h
P0H	b FF02h	81h	Port 0 High Register (Upper half of PORT0)	00h
P1L	b FF04h	82h	Port 1 Low Register (Lower half of PORT1)	00h
P1H	b FF06h	83h	Port 1 High Register (Upper half of PORT1)	00h
P2	b FFC0h	E0h	Port 2 Register	0000h
P3	b FFC4h	E2h	Port 3 Register	0000h
P4	b FFC8h	E4h	Port 4 Register (8 bits)	00h
P5	b FFA2h	D1h	Port 5 Register (read only)	XXXXh
P6	b FFCh	E6h	Port 6 Register (8 bits)	00h
PECC0	FEC0h	60h	PEC Channel 0 Control Register	0000h

## Special Function Register Overview (Cont'd)

Name	Physical Address	8-Bit Address	Description	Reset Value
PECC1	FEC2h	61h	PEC Channel 1 Control Register	0000h
PECC2	FEC4h	62h	PEC Channel 2 Control Register	0000h
PECC3	FEC6h	63h	PEC Channel 3 Control Register	0000h
PECC4	FEC8h	64h	PEC Channel 4 Control Register	0000h
PECC5	FECAh	65h	PEC Channel 5 Control Register	0000h
PECC6	FECCh	66h	PEC Channel 6 Control Register	0000h
PECC7	FECEh	67h	PEC Channel 7 Control Register	0000h
PSW	b FF10h	88h	CPU Program Status Word	0000h
RP0H	b F108h E	84h	System Start-up Configuration Register (Rd. only)	XXh
S0BG	FEB4h	5Ah	Serial Channel 0 Baud Rate Generator Reload Register	0000h
S0CON	b FFB0h	D8h	Serial Channel 0 Control Register	0000h
S0EIC	b FF70h	B8h	Serial Channel 0 Error Interrupt Control Register	0000h
S0RBUF	FEB2h	59h	Serial Channel 0 Receive Buffer Register (read only)	XXh
S0RIC	b FF6Eh	B7h	Serial Channel 0 Receive Interrupt Control Register	0000h
S0TBIC	b F19Ch E	CEh	Serial Channel 0 Transmit Buffer Interrupt Control Register	0000h
S0TBUF	FEB0h	58h	Serial Channel 0 Transmit Buffer Register (write only)	00h
S0TIC	b FF6Ch	B6h	Serial Channel 0 Transmit Interrupt Control Register	0000h
SP	FE12h	09h	CPU System Stack Pointer Register	FC00h
SSCBR	F0B4h E	5Ah	SSC Baudrate Register	0000h
SSCCON	b FFB2h	D9h	SSC Control Register	0000h
SSCEIC	b FF76h	BBh	SSC Error Interrupt Control Register	0000h
SSCRB	F0B2h E	59h	SSC Receive Buffer (read only)	XXXXh
SSCRIC	b FF74h	BAh	SSC Receive Interrupt Control Register	0000h
SSCTB	F0B0h E	58h	SSC Transmit Buffer (write only)	0000h
SSCTIC	b FF72h	B9h	SSC Transmit Interrupt Control Register	0000h
STKOV	FE14h	0Ah	CPU Stack Overflow Pointer Register	FA00h
STKUN	FE16h	0Bh	CPU Stack Underflow Pointer Register	FC00h
SYSCON	b FF12h	89h	CPU System Configuration Register	0xx0h*)
T2	FE40h	20h	GPT1 Timer 2 Register	0000h
T2CON	b FF40h	A0h	GPT1 Timer 2 Control Register	0000h

### Special Function Register Overview (Cont'd)

Name	Physical Address	8-Bit Address	Description	Reset Value
T2IC	b FF60h	B0h	GPT1 Timer 2 Interrupt Control Register	0000h
T3	FE42h	21h	GPT1 Timer 3 Register	0000h
T3CON	b FF42h	A1h	GPT1 Timer 3 Control Register	0000h
T3IC	b FF62h	B1h	GPT1 Timer 3 Interrupt Control Register	0000h
T4	FE44h	22h	GPT1 Timer 4 Register	0000h
T4CON	b FF44h	A2h	GPT1 Timer 4 Control Register	0000h
T4IC	b FF64h	B2h	GPT1 Timer 4 Interrupt Control Register	0000h
T5	FE46h	23h	GPT2 Timer 5 Register	0000h
T5CON	b FF46h	A3h	GPT2 Timer 5 Control Register	0000h
T5IC	b FF66h	B3h	GPT2 Timer 5 Interrupt Control Register	0000h
T6	FE48h	24h	GPT2 Timer 6 Register	0000h
T6CON	b FF48h	A4h	GPT2 Timer 6 Control Register	0000h
T6IC	b FF68h	B4h	GPT2 Timer 6 Interrupt Control Register	0000h
TFR	b FFACh	D6h	Trap Flag Register	0000h
WDT	FEAEh	57h	Watchdog Timer Register (read only)	0000h
WDTCON	FFAEh	D7h	Watchdog Timer Control Register	0000h
XP0IC	b F186h <sub>E</sub>	C3h	X-Peripheral 0 Interrupt Control Register	0000h
XP1IC	b F18Eh <sub>E</sub>	C7h	X-Peripheral 1 Interrupt Control Register	0000h
XP2IC	b F196h <sub>E</sub>	CBh	X-Peripheral 2 Interrupt Control Register	0000h
XP3IC	b F19Eh <sub>E</sub>	CFh	X-Peripheral 3 Interrupt Control Register	0000h
ZEROS	b FF1Ch	8Eh	Constant Value 0's Register (read only)	0000h

\*) The system configuration is selected during reset.

**Note:** The Interrupt Control Registers XPnIC are prepared to control interrupt requests from integrated X-Bus peripherals. Nodes, where no X-Peripherals are connected, may be used to generate software controlled interrupt requests by setting the respective XPnIR bit.

## 13 ELECTRICAL CHARACTERISTICS

### 13.1 Absolute Maximum Ratings

- Ambient temperature under bias ( $T_A$ ): 0 to +70 °C
- Storage temperature ( $T_{ST}$ ): – 65 to +150 °C
- Voltage on  $V_{DD}$  pins with respect to ground ( $V_{SS}$ ): 0.5 to +6.5 V
- Voltage on any pin with respect to ground ( $V_{SS}$ ): –0.5 to  $V_{DD} + 0.5$  V
- Input current on any pin during overload condition: –10 to +10 mA
- Absolute sum of all input currents during overload condition: |100 mA|
- Power dissipation: 1.5 W

**Note:** Stresses above those listed under “Absolute Maximum Ratings” may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability. During overload conditions ( $V_{IN} > V_{DD}$  or  $V_{IN} < V_{SS}$ ) the voltage on pins with respect to ground ( $V_{SS}$ ) must not exceed the values defined by the Absolute Maximum Ratings.

### 13.2 Parameter Interpretation

The parameters listed in the following partly represent the characteristics of the ST10R165 and partly its demands on the system. To aid in interpreting the parameters right, when evaluating them for a design, they are marked in column “Symbol”:

**CC (Controller Characteristics):**

The logic of the ST10R165 will provide signals with the respective timing characteristics.

**SR (System Requirement):**

The external system must provide signals with the respective timing characteristics to the ST10R165.

## 13.3 DC Characteristics

 $V_{DD} = 5\text{ V} \pm 10\%$ ;  $V_{SS} = 0\text{ V}$ ;  $f_{CPU} = 20\text{ MHz}$ ; Reset active.  $T_A = 0\text{ to }+70\text{ }^\circ\text{C}$ 

Parameter	Symbol	Limit Values		Unit	Test Condition
		min.	max.		
Input low voltage	$V_{IL}$ SR	-0.5	$0.2 V_{DD} - 0.1$	V	-
Input high voltage (all except $\overline{RSTIN}$ and XTAL1)	$V_{IH}$ SR	$0.2 V_{DD} + 0.9$	$V_{DD} + 0.5$	V	-
Input high voltage $\overline{RSTIN}$	$V_{IH1}$ SR	$0.6 V_{DD}$	$V_{DD} + 0.5$	V	-
Input high voltage XTAL1	$V_{IH2}$ SR	$0.7 V_{DD}$	$V_{DD} + 0.5$	V	-
Output low voltage (PORT0, PORT1, Port 4, ALE, $\overline{RD}$ , $\overline{WR}$ , $\overline{BHE}$ , CLKOUT, $\overline{RSTOUT}$ )	$V_{OL}$ CC	-	0.45	V	$I_{OL} = 2.4\text{ mA}$
Output low voltage (all other outputs)	$V_{OL1}$ CC	-	0.45	V	$I_{OL1} = 1.6\text{ mA}$
Output high voltage (PORT0, PORT1, Port 4, ALE, $\overline{RD}$ , $\overline{WR}$ , $\overline{BHE}$ , CLKOUT, $\overline{RSTOUT}$ )	$V_{OH}$ CC	$0.9 V_{DD}$ 2.4	-	V	$I_{OH} = -500\text{ }\mu\text{A}$ $I_{OH} = -2.4\text{ mA}$
Output high voltage <sup>1)</sup> (all other outputs)	$V_{OH1}$ CC	$0.9 V_{DD}$ 2.4	-	V V	$I_{OH} = -250\text{ }\mu\text{A}$ $I_{OH} = -1.6\text{ mA}$
Input leakage current (Port 5)	$I_{OZ1}$ CC	-	$\pm 200$	nA	$0\text{ V} < V_{IN} < V_{DD}$
Input leakage current (all other)	$I_{OZ2}$ CC	-	$\pm 500$	nA	$0\text{ V} < V_{IN} < V_{DD}$
$\overline{RSTIN}$ pull-up resistor	$R_{RST}$ CC	50	150	k $\Omega$	-
Read/Write inactive current <sup>4)</sup>	$I_{RWH}$ <sup>2)</sup>	-	-40	$\mu\text{A}$	$V_{OUT} = 2.4\text{ V}$
Read/Write active current <sup>4)</sup>	$I_{RWL}$ <sup>3)</sup>	-500	-	$\mu\text{A}$	$V_{OUT} = V_{OLmax}$
ALE inactive current <sup>4)</sup>	$I_{ALEL}$ <sup>2)</sup>	-	40	$\mu\text{A}$	$V_{OUT} = V_{OLmax}$
ALE active current <sup>4)</sup>	$I_{ALEH}$ <sup>3)</sup>	500	-	$\mu\text{A}$	$V_{OUT} = 2.4\text{ V}$
Port 6 inactive current <sup>4)</sup>	$I_{P6H}$ <sup>2)</sup>	-	-40	$\mu\text{A}$	$V_{OUT} = 2.4\text{ V}$
Port 6 active current <sup>4)</sup>	$I_{P6L}$ <sup>3)</sup>	-500	-	$\mu\text{A}$	$V_{OUT} = V_{OL1max}$

## DC Characteristics (Cont'd)

 $V_{DD} = 5\text{ V} \pm 10\%$ ;  $V_{SS} = 0\text{ V}$ ;  $f_{CPU} = 20\text{ MHz}$ ; Reset active.  $T_A = 0\text{ to }+70\text{ }^\circ\text{C}$ 

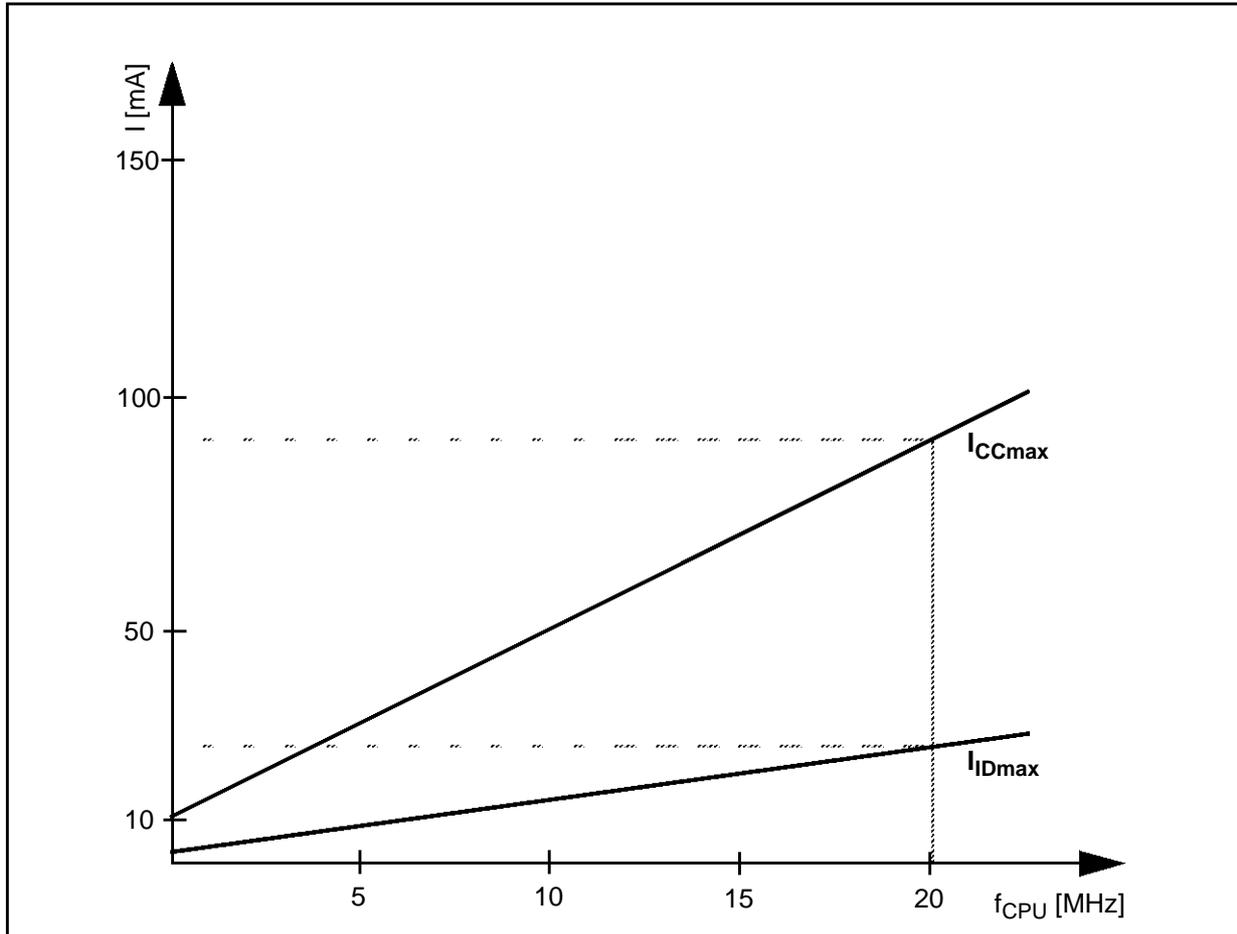
Parameter	Symbol	Limit Values		Unit	Test Condition
		min.	max.		
PORT0 configuration current <sup>4)</sup>	$I_{POH}$ <sup>2)</sup>	–	-10	$\mu\text{A}$	$V_{IN} = V_{IHmin}$
	$I_{POL}$ <sup>3)</sup>	-100	–	$\mu\text{A}$	$V_{IN} = V_{ILmax}$
XTAL1 input current	$I_{IL}$ CC	–	$\pm 20$	$\mu\text{A}$	$0\text{ V} < V_{IN} < V_{DD}$
Pin capacitance <sup>5)</sup> (digital inputs/outputs)	$C_{IO}$ CC	–	10	$\text{pF}$	$f = 1\text{ MHz}$ $T_A = 25\text{ }^\circ\text{C}$
Power supply current	$I_{CC}$	–	$10 + 4 * f_{CPU}$	$\text{mA}$	$\overline{RSTIN} = V_{IL2}$ $f_{CPU}$ in [MHz] <sup>6)</sup>
Idle mode supply current	$I_{ID}$	–	$2 + 1.2 * f_{CPU}$	$\text{mA}$	$\overline{RSTIN} = V_{IH1}$ $f_{CPU}$ in [MHz] <sup>6)</sup>
Power-down mode supply current	$I_{PD}$	–	100	$\mu\text{A}$	$V_{DD} = 5.5\text{ V}$ <sup>7)</sup>

## Notes:

- 1) This specification is not valid for outputs which are switched to open drain mode. In this case the respective output will float and the voltage results from the external circuitry.
- 2) The maximum current may be drawn while the respective signal line remains inactive.
- 3) The minimum current must be drawn in order to drive the respective signal line active.
- 4) This specification is only valid during Reset, or during Hold- or Adapt-mode. Port 6 pins are only affected, if they are used for  $\overline{CS}$  output and the open drain function is not enabled.
- 5) Not 100% tested, guaranteed by design characterization.
- 6) The supply current is a function of the operating frequency. This dependency is illustrated in the figure below.  
These parameters are tested at  $V_{DDmax}$  and 20 MHz CPU clock with all outputs disconnected and all inputs at  $V_{IL}$  or  $V_{IH}$ .
- 7) This parameter is tested including leakage currents. All inputs (including pins configured as inputs) at 0 V to 0.1 V or at  $V_{DD} - 0.1\text{ V}$  to  $V_{DD}$ ,  $V_{REF} = 0\text{ V}$ , all outputs (including pins configured as outputs) disconnected.

## DC Characteristics (Cont'd)

Figure 8. Supply/Idle Current as a Function of Operating Frequency



13.4 Testing Waveforms

Figure 9. Input Output Waveforms

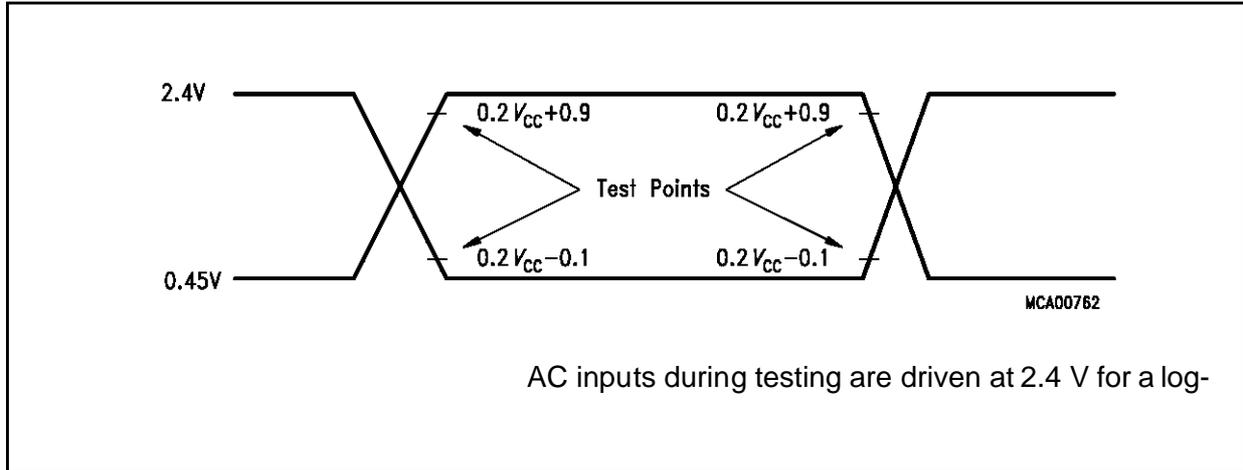
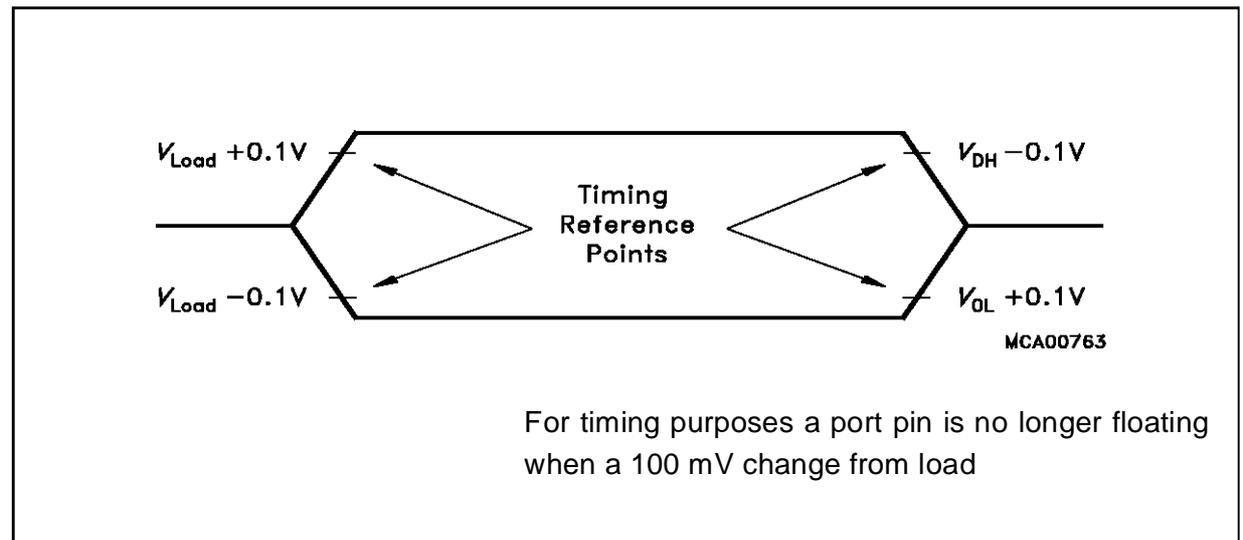


Figure 10. Float Waveforms



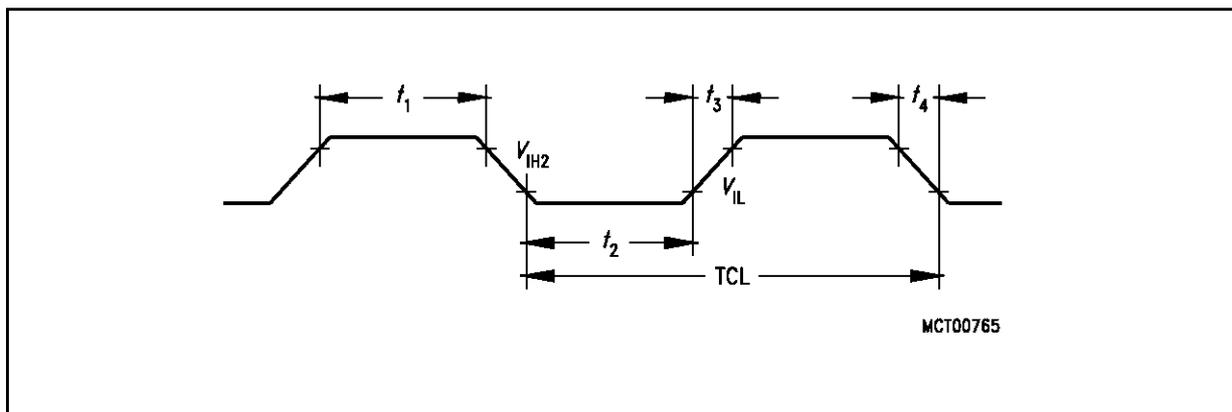
### 13.5 AC Characteristics

#### External Clock Drive XTAL1

$$V_{DD} = 5\text{ V} \pm 10\%; \quad V_{SS} = 0\text{ V} \quad T_A = 0\text{ to }+70\text{ }^\circ\text{C}$$

Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
Oscillator period	TCL SR	25	25	25	500	ns
High time	$t_1$ SR	6	–	6	–	ns
Low time	$t_2$ SR	6	–	6	–	ns
Rise time	$t_3$ SR	–	5	–	5	ns
Fall time	$t_4$ SR	–	5	–	5	ns

Figure 11. External Clock Drive XTAL1



#### 13.5.1 Memory Cycle Variables

The timing tables below use three variables which are derived from the BUSCONx registers and represent the special characteristics of the programmed memory cycle. The following table describes, how these variables are to be computed.

Description	Symbol	Values
ALE Extension	$t_A$	$TCL * \langle ALECTL \rangle$
Memory Cycle Time Waitstates	$t_C$	$2TCL * (15 - \langle MCTC \rangle)$
Memory Tristate Time	$t_F$	$2TCL * (1 - \langle MTTC \rangle)$

## 13.5.2 Multiplexed Bus

$$V_{DD} = 5\text{ V} \pm 10\%; V_{SS} = 0\text{ V}$$

$$T_A = 0\text{ to }+70\text{ }^\circ\text{C}$$

$$C_L \text{ (for PORT0, PORT1, Port 4, ALE, } \overline{\text{RD}}, \overline{\text{WR}}, \overline{\text{BHE}}, \text{CLKOUT)} = 100\text{ pF}$$

$$C_L \text{ (for Port 6, } \overline{\text{CS}}) = 100\text{ pF}$$

$$\text{ALE cycle time} = 6\text{ TCL} + 2t_A + t_C + t_F \text{ (150 ns at 20-MHz CPU clock without waitstates)}$$

Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
ALE high time	$t_5$ CC	$15 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
Address setup to ALE	$t_6$ CC	$10 + t_A$	–	$\text{TCL} - 15 + t_A$	–	ns
Address hold after ALE	$t_7$ CC	$15 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
ALE falling edge to $\overline{\text{RD}}$ , $\overline{\text{WR}}$ (with RW-delay)	$t_8$ CC	$15 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
ALE falling edge to $\overline{\text{RD}}$ , $\overline{\text{WR}}$ (no RW-delay)	$t_9$ CC	$-10 + t_A$	–	$-10 + t_A$	–	ns
Address float after $\overline{\text{RD}}$ , $\overline{\text{WR}}$ (with RW-delay)	$t_{10}$ CC	–	5	–	5	ns
Address float after $\overline{\text{RD}}$ , $\overline{\text{WR}}$ (no RW-delay)	$t_{11}$ CC	–	30	–	$\text{TCL} + 5$	ns
$\overline{\text{RD}}$ , $\overline{\text{WR}}$ low time (with RW-delay)	$t_{12}$ CC	$40 + t_C$	–	$2\text{TCL} - 10 + t_C$	–	ns
$\overline{\text{RD}}$ , $\overline{\text{WR}}$ low time (no RW-delay)	$t_{13}$ CC	$65 + t_C$	–	$3\text{TCL} - 10 + t_C$	–	ns
$\overline{\text{RD}}$ to valid data in (with RW-delay)	$t_{14}$ SR	–	$30 + t_C$	–	$2\text{TCL} - 20 + t_C$	ns
$\overline{\text{RD}}$ to valid data in (no RW-delay)	$t_{15}$ SR	–	$55 + t_C$	–	$3\text{TCL} - 20 + t_C$	ns
ALE low to valid data in	$t_{16}$ SR	–	$55 + t_A + t_C$	–	$3\text{TCL} - 20 + t_A + t_C$	ns
Address to valid data in	$t_{17}$ SR	–	$70 + 2t_A + t_C$	–	$4\text{TCL} - 30 + 2t_A + t_C$	ns
Data hold after $\overline{\text{RD}}$ rising edge	$t_{18}$ SR	0	–	0	–	ns
Data float after $\overline{\text{RD}}$	$t_{19}$ SR	–	$35 + t_F$	–	$2\text{TCL} - 15 + t_F$	ns
Data valid to $\overline{\text{WR}}$	$t_{22}$ SR	$35 + t_C$	–	$2\text{TCL} - 15 + t_C$	–	ns
Data hold after $\overline{\text{WR}}$	$t_{23}$ CC	$35 + t_F$	–	$2\text{TCL} - 15 + t_F$	–	ns
ALE rising edge after $\overline{\text{RD}}$ , $\overline{\text{WR}}$	$t_{25}$ CC	$35 + t_F$	–	$2\text{TCL} - 15 + t_F$	–	ns

## Multiplexed Bus (Cont'd)

Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
Address hold after $\overline{RD}$ , $\overline{WR}$	$t_{27}$ CC	$35 + t_F$	–	$2TCL - 15 + t_F$	–	ns
ALE falling edge to $\overline{CS}$	$t_{38}$ CC	$-5 - t_A$	$10 - t_A$	$-5 - t_A$	$10 - t_A$	ns
$\overline{CS}$ low to Valid Data In	$t_{39}$ SR	–	$55 + t_C + 2t_A$	–	$3TCL - 20 + t_C + 2t_A$	ns
$\overline{CS}$ hold after $\overline{RD}$ , $\overline{WR}$	$t_{40}$ CC	$60 + t_F$	–	$3TCL - 15 + t_F$	–	ns
ALE fall. edge to $\overline{RdCS}$ , $\overline{WrCS}$ (with RW delay)	$t_{42}$ CC	$20 + t_A$	–	$TCL - 5 + t_A$	–	ns
ALE fall. edge to $\overline{RdCS}$ , $\overline{WrCS}$ (no RW delay)	$t_{43}$ CC	$-5 + t_A$	–	$-5 + t_A$	–	ns
Address float after $\overline{RdCS}$ , $\overline{WrCS}$ (with RW delay)	$t_{44}$ CC	–	0	–	0	ns
Address float after $\overline{RdCS}$ , $\overline{WrCS}$ (no RW delay)	$t_{45}$ CC	–	25	–	TCL	ns
$\overline{RdCS}$ to Valid Data In (with RW delay)	$t_{46}$ SR	–	$25 + t_C$	–	$2TCL - 25 + t_C$	ns
$\overline{RdCS}$ to Valid Data In (no RW delay)	$t_{47}$ SR	–	$50 + t_C$	–	$3TCL - 25 + t_C$	ns
$\overline{RdCS}$ , $\overline{WrCS}$ Low Time (with RW delay)	$t_{48}$ CC	$40 + t_C$	–	$2TCL - 10 + t_C$	–	ns
$\overline{RdCS}$ , $\overline{WrCS}$ Low Time (no RW delay)	$t_{49}$ CC	$65 + t_C$	–	$3TCL - 10 + t_C$	–	ns
Data valid to $\overline{WrCS}$	$t_{50}$ CC	$35 + t_C$	–	$2TCL - 15 + t_C$	–	ns
Data hold after $\overline{RdCS}$	$t_{51}$ SR	0	–	0	–	ns
Data float after $\overline{RdCS}$	$t_{52}$ SR	–	$30 + t_F$	–	$2TCL - 20 + t_F$	ns
Address hold after $\overline{RdCS}$ , $\overline{WrCS}$	$t_{54}$ CC	$30 + t_F$	–	$2TCL - 20 + t_F$	–	ns
Data hold after $\overline{WrCS}$	$t_{56}$ CC	$30 + t_F$	–	$2TCL - 20 + t_F$	–	ns

Figure 12. External Memory Cycle: Multiplexed Bus, With Read/Write Delay, Normal ALE

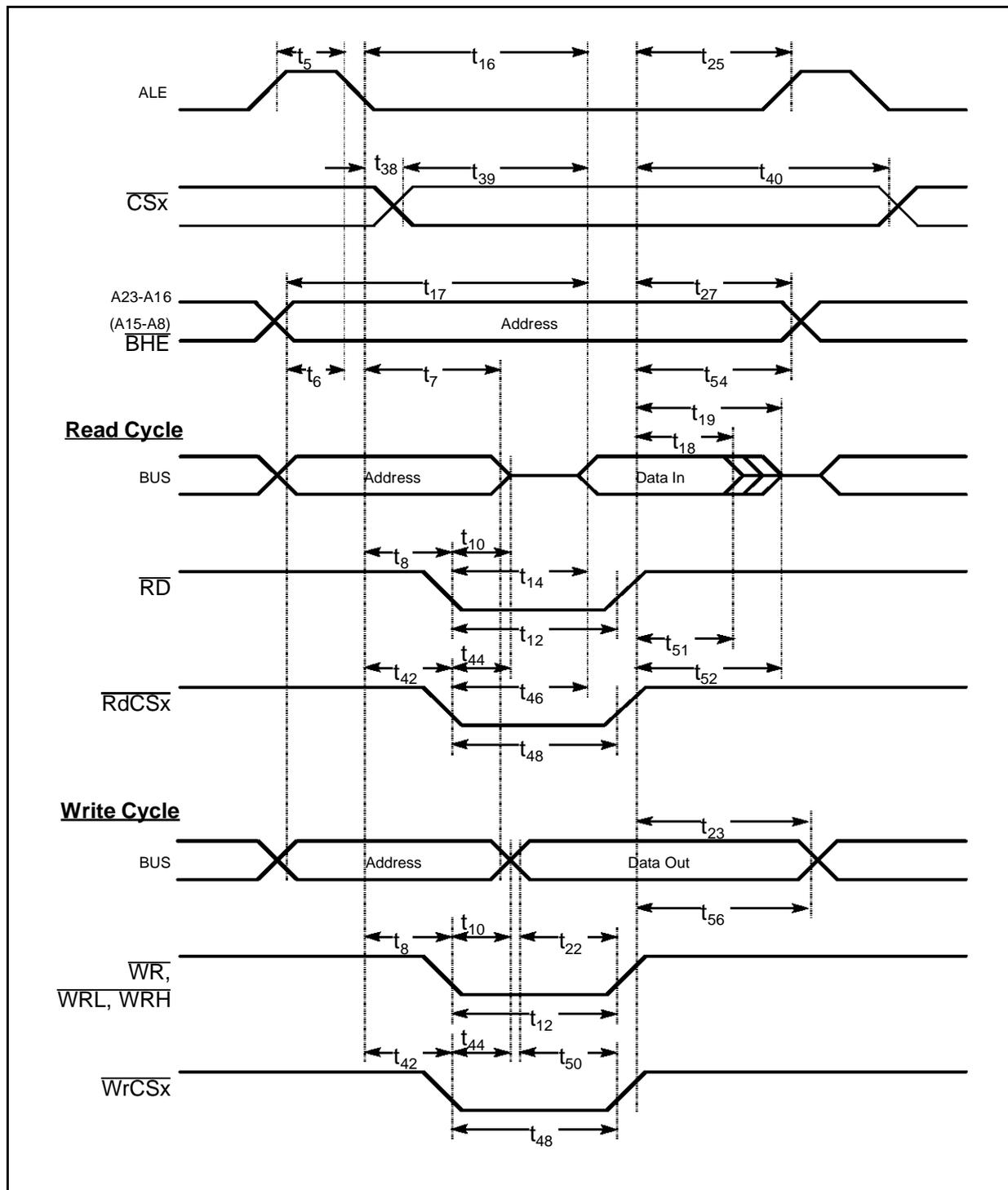


Figure 13. External Memory Cycle: Multiplexed Bus, With Read/Write Delay, Extended ALE

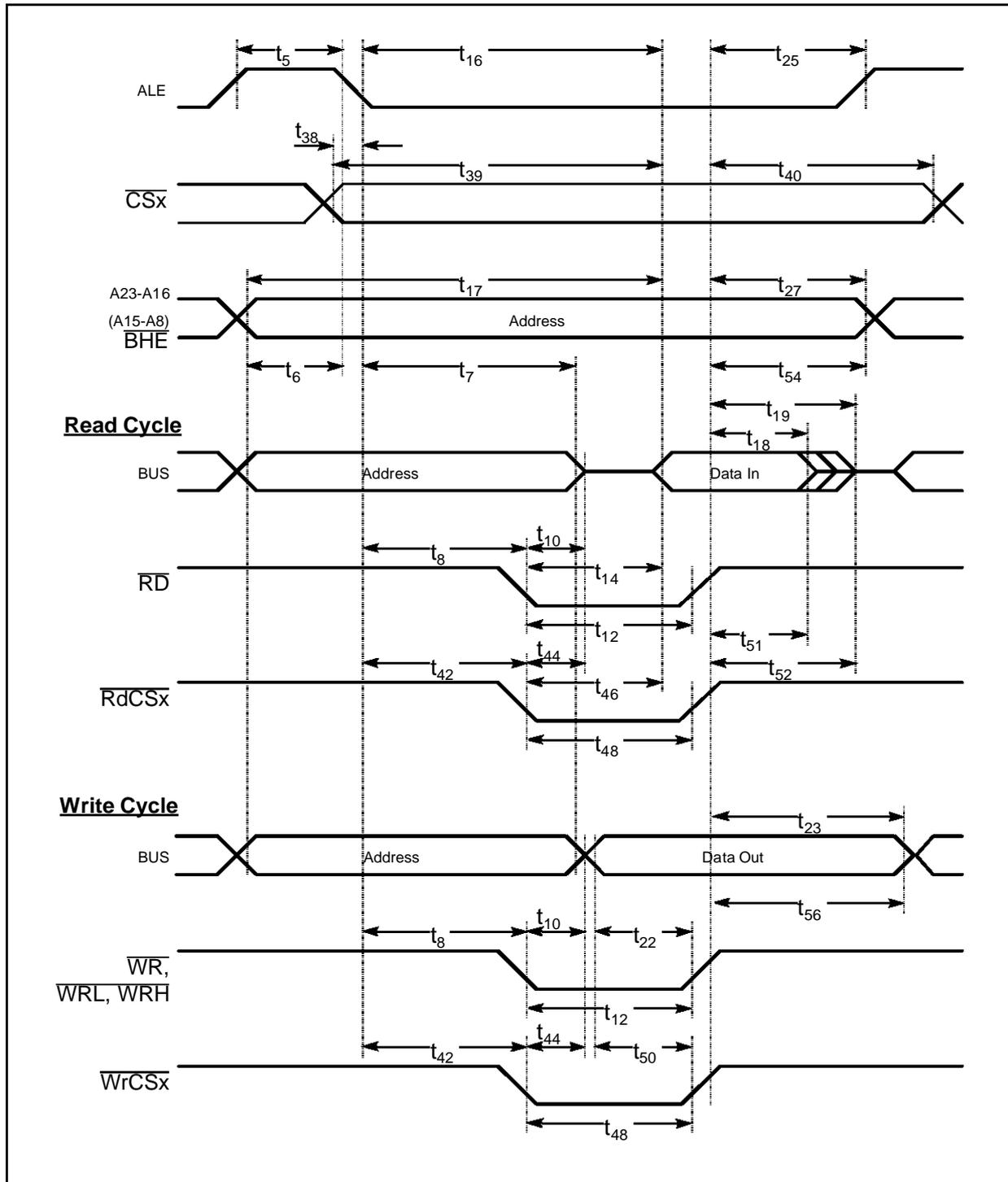


Figure 14. External Memory Cycle: Multiplexed Bus, No Read/Write Delay, Normal ALE

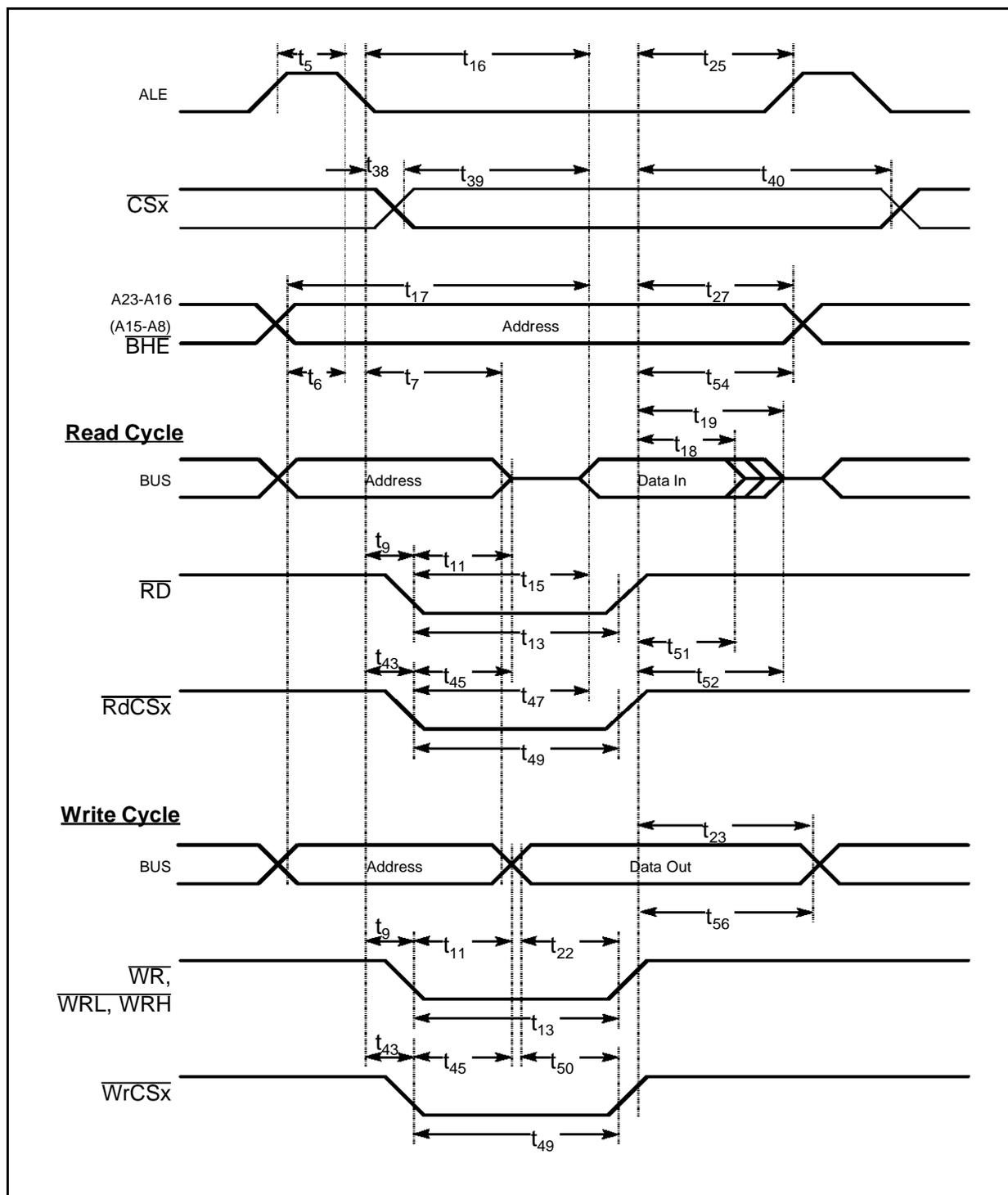
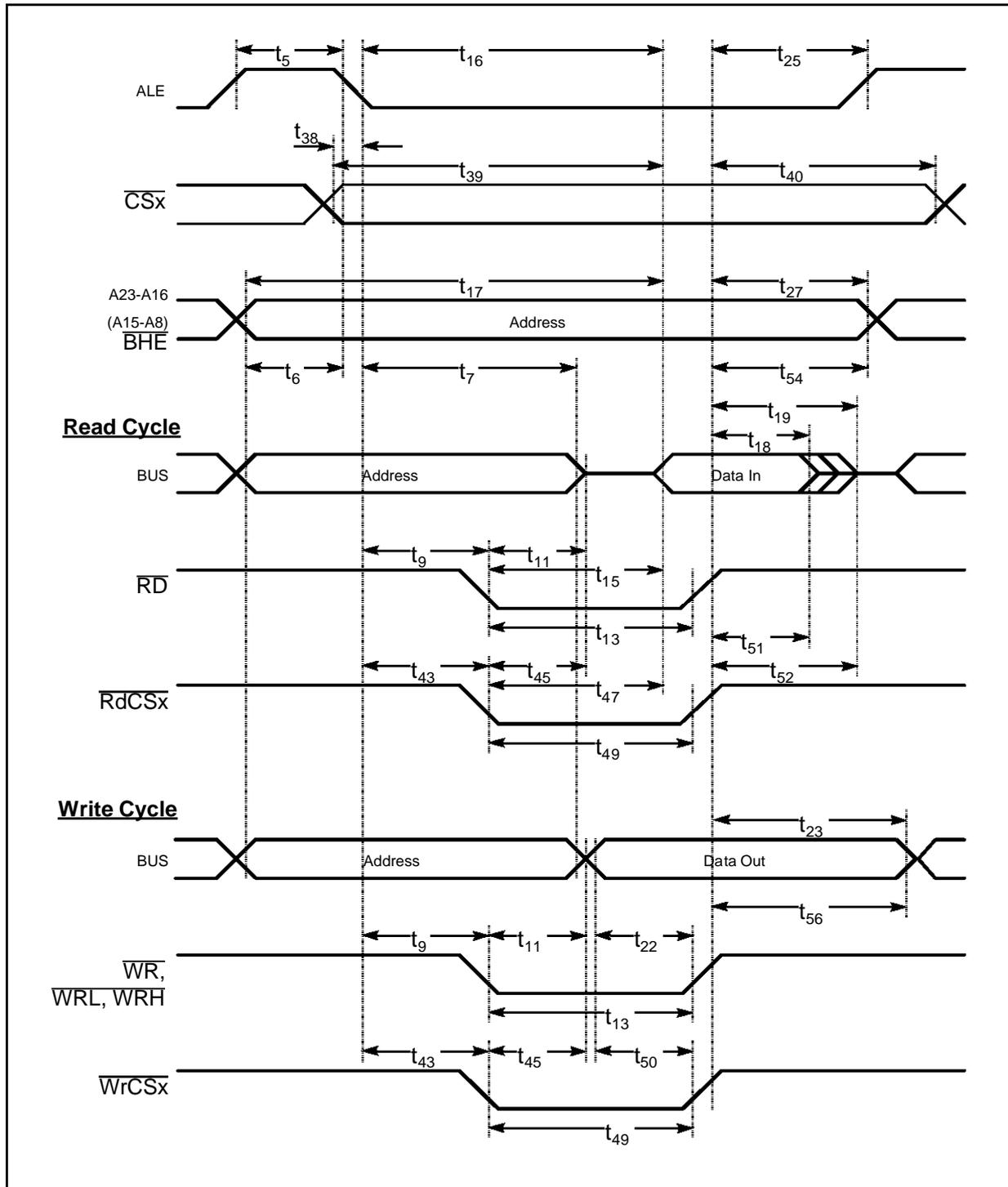


Figure 15. External Memory Cycle: Multiplexed Bus, No Read/Write Delay, Extended ALE



## 13.5.2.1 Demultiplexed Bus

$$V_{DD} = 5\text{ V} \pm 10\%; V_{SS} = 0\text{ V}$$

$$T_A = 0\text{ to }+70\text{ }^\circ\text{C}$$

$$C_L \text{ (for PORT0, PORT1, Port 4, ALE, } \overline{\text{RD}}, \overline{\text{WR}}, \overline{\text{BHE}}, \text{CLKOUT)} = 100\text{ pF}$$

$$C_L \text{ (for Port 6, } \overline{\text{CS}}) = 100\text{ pF}$$

$$\text{ALE cycle time} = 4\text{ TCL} + 2t_A + t_C + t_F \text{ (100 ns at 20 MHz CPU clock without waitstates)}$$

Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
ALE high time	$t_5$ CC	$15 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
Address setup to ALE	$t_6$ CC	$10 + t_A$	–	$\text{TCL} - 15 + t_A$	–	ns
ALE falling edge to $\overline{\text{RD}}$ , $\overline{\text{WR}}$ (with RW-delay)	$t_8$ CC	$15 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
ALE falling edge to $\overline{\text{RD}}$ , $\overline{\text{WR}}$ (no RW-delay)	$t_9$ CC	$-10 + t_A$	–	$-10 + t_A$	–	ns
$\overline{\text{RD}}$ , $\overline{\text{WR}}$ low time (with RW-delay)	$t_{12}$ CC	$40 + t_C$	–	$2\text{TCL} - 10 + t_C$	–	ns
$\overline{\text{RD}}$ , $\overline{\text{WR}}$ low time (no RW-delay)	$t_{13}$ CC	$65 + t_C$	–	$3\text{TCL} - 10 + t_C$	–	ns
$\overline{\text{RD}}$ to valid data in (with RW-delay)	$t_{14}$ SR	–	$30 + t_C$	–	$2\text{TCL} - 20 + t_C$	ns
$\overline{\text{RD}}$ to valid data in (no RW-delay)	$t_{15}$ SR	–	$55 + t_C$	–	$3\text{TCL} - 20 + t_C$	ns
ALE low to valid data in	$t_{16}$ SR	–	$55 + t_A + t_C$	–	$3\text{TCL} - 20 + t_A + t_C$	ns
Address to valid data in	$t_{17}$ SR	–	$70 + 2t_A + t_C$	–	$4\text{TCL} - 30 + 2t_A + t_C$	ns
Data hold after $\overline{\text{RD}}$ rising edge	$t_{18}$ SR	0	–	0	–	ns
Data float after $\overline{\text{RD}}$ rising edge (with RW-delay)	$t_{20}$ SR	–	$35 + t_F$	–	$2\text{TCL} - 15 + t_F$	ns
Data float after $\overline{\text{RD}}$ rising edge (no RW-delay)	$t_{21}$ SR	–	$15 + t_F$	–	$\text{TCL} - 10 + t_F$	ns
Data valid to $\overline{\text{WR}}$	$t_{22}$ CC	$35 + t_C$	–	$2\text{TCL} - 15 + t_C$	–	ns
Data hold after $\overline{\text{WR}}$	$t_{24}$ CC	$15 + t_F$	–	$\text{TCL} - 10 + t_F$	–	ns
ALE rising edge after $\overline{\text{RD}}$ , $\overline{\text{WR}}$	$t_{26}$ CC	$-10 + t_F$	–	$-10 + t_F$	–	ns

## Demultiplexed Bus (Cont'd)

Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
Address hold after $\overline{RD}$ , $\overline{WR}$	$t_{28}$ CC	$0 + t_F$	–	$0 + t_F$	–	ns
ALE falling edge to $\overline{CS}$	$t_{38}$ CC	$-5 - t_A$	$10 - t_A$	$-5 - t_A$	$10 - t_A$	ns
$\overline{CS}$ low to Valid Data In	$t_{39}$ SR	–	$55 + t_C + 2t_A$	–	$3TCL - 20 + t_C + 2t_A$	ns
$\overline{CS}$ hold after $\overline{RD}$ , $\overline{WR}$	$t_{41}$ CC	$10 + t_F$	–	$TCL - 15 + t_F$	–	ns
ALE falling edge to $\overline{RdCS}$ , $\overline{WrCS}$ (with RW-delay)	$t_{42}$ CC	$20 + t_A$	–	$TCL - 5 + t_A$	–	ns
ALE falling edge to $\overline{RdCS}$ , $\overline{WrCS}$ (no RW-delay)	$t_{43}$ CC	$-5 + t_A$	–	$-5 + t_A$	–	ns
$\overline{RdCS}$ to Valid Data In (with RW-delay)	$t_{46}$ SR	–	$25 + t_C$	–	$2TCL - 25 + t_C$	ns
$\overline{RdCS}$ to Valid Data In (no RW-delay)	$t_{47}$ SR	–	$50 + t_C$	–	$3TCL - 25 + t_C$	ns
$\overline{RdCS}$ , $\overline{WrCS}$ Low Time (with RW-delay)	$t_{48}$ CC	$40 + t_C$	–	$2TCL - 10 + t_C$	–	ns
$\overline{RdCS}$ , $\overline{WrCS}$ Low Time (no RW-delay)	$t_{49}$ CC	$65 + t_C$	–	$3TCL - 10 + t_C$	–	ns
Data valid to $\overline{WrCS}$	$t_{50}$ CC	$35 + t_C$	–	$2TCL - 15 + t_C$	–	ns
Data hold after $\overline{RdCS}$	$t_{51}$ SR	0	–	0	–	ns
Data float after $\overline{RdCS}$ (with RW-delay)	$t_{53}$ SR	–	$30 + t_F$	–	$2TCL - 20 + t_F$	ns
Data float after $\overline{RdCS}$ (no RW-delay)	$t_{68}$ SR	–	$5 + t_F$	–	$TCL - 20 + t_F$	ns
Address hold after $\overline{RdCS}$ , $\overline{WrCS}$	$t_{55}$ CC	$-5 + t_F$	–	$-5 + t_F$	–	ns
Data hold after $\overline{WrCS}$	$t_{57}$ CC	$10 + t_F$	–	$TCL - 15 + t_F$	–	ns

Figure 16. External Memory Cycle: Demultiplexed Bus, With Read/Write Delay, Normal ALE

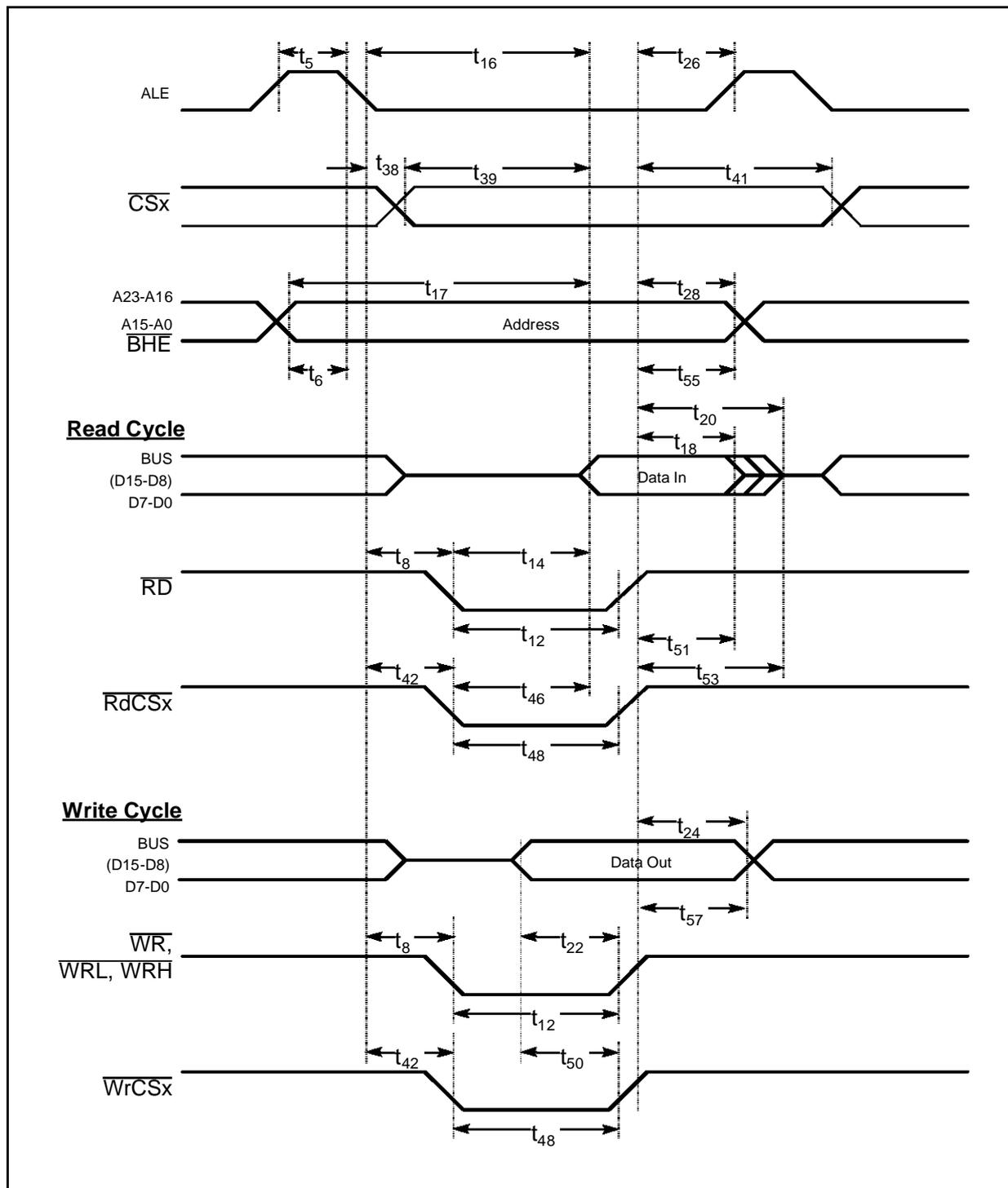


Figure 17. External Memory Cycle: Demultiplexed Bus, With Read/Write Delay, Extended ALE

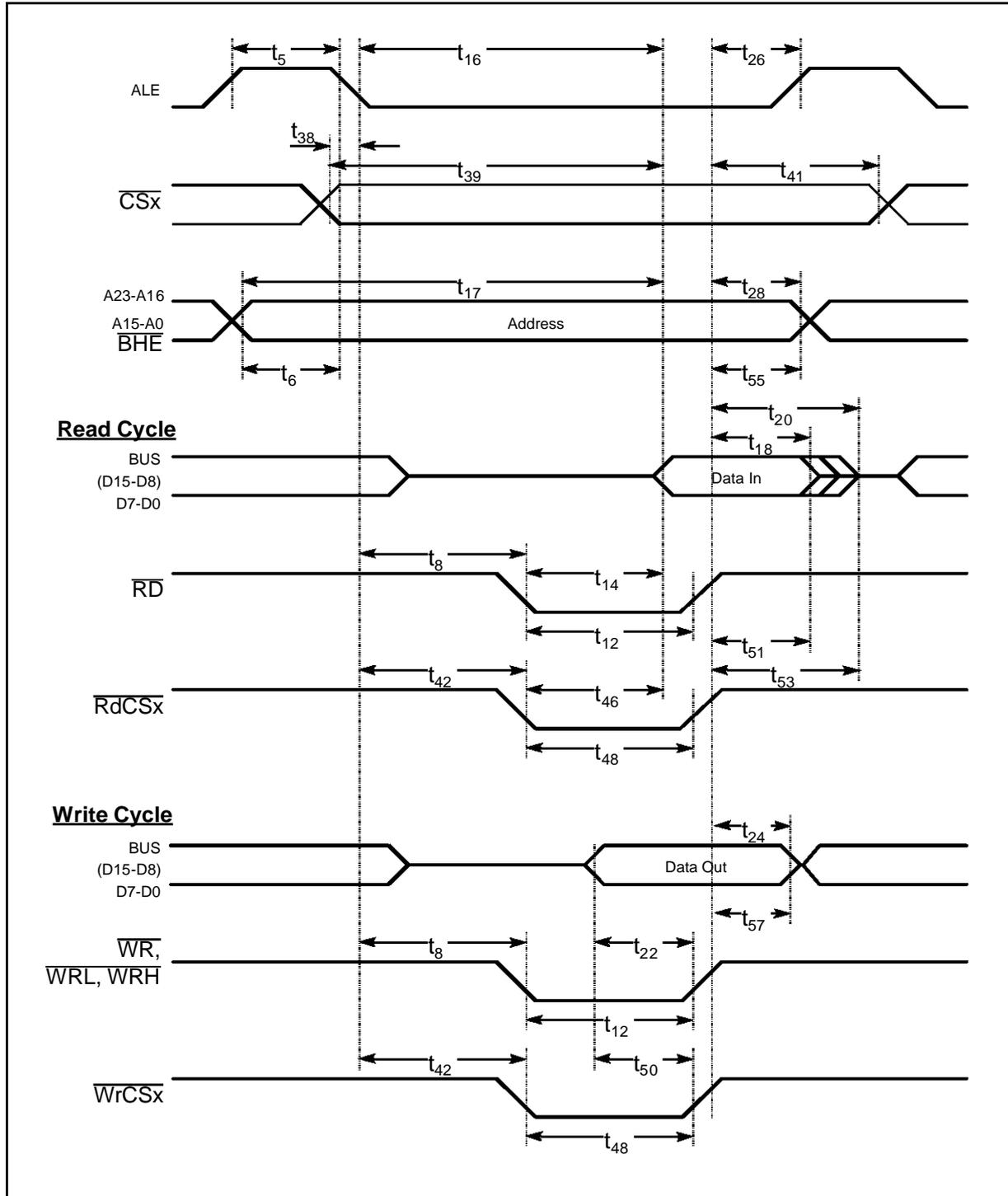


Figure 18. External Memory Cycle: Demultiplexed Bus, No Read/Write Delay, Normal ALE

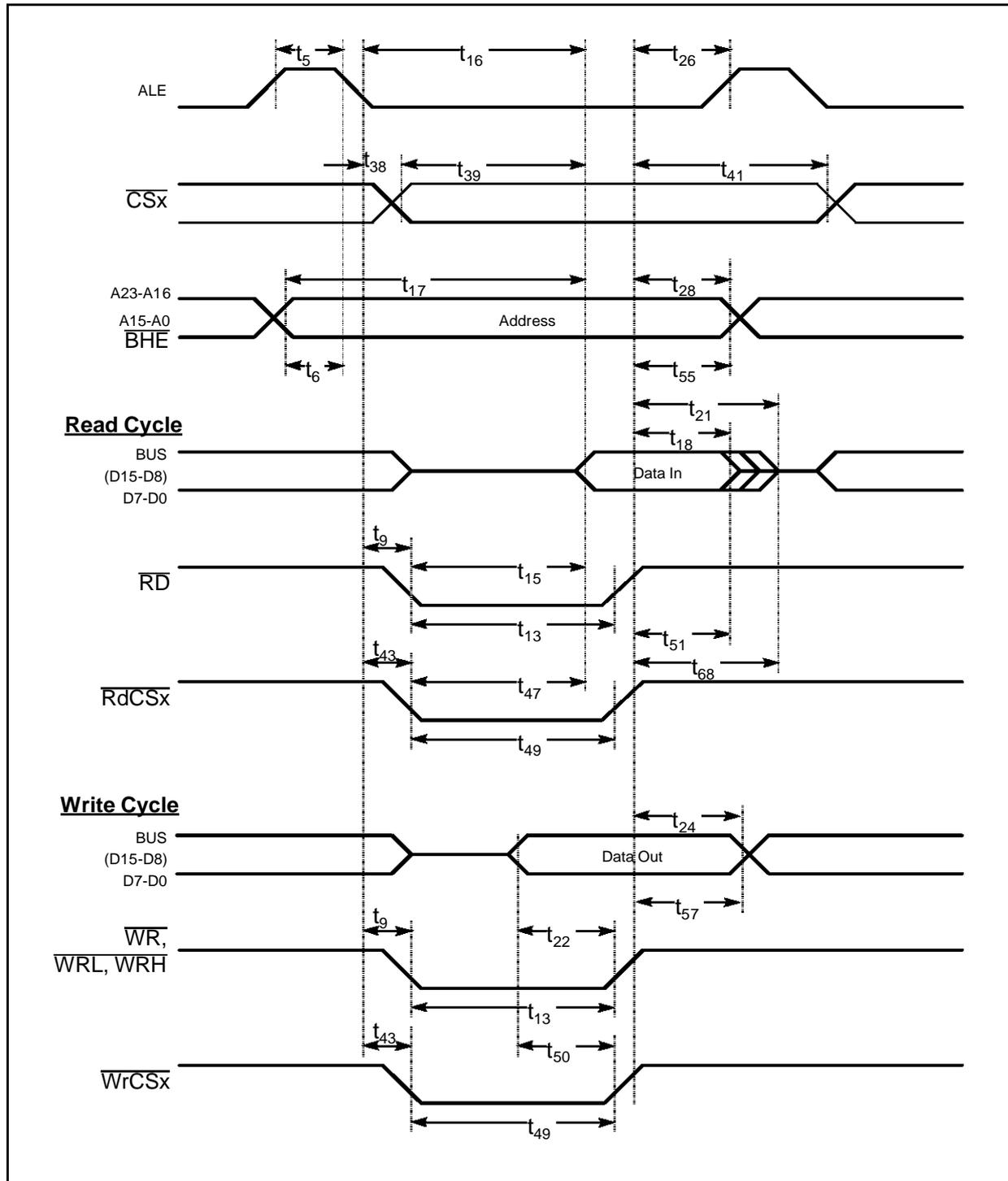
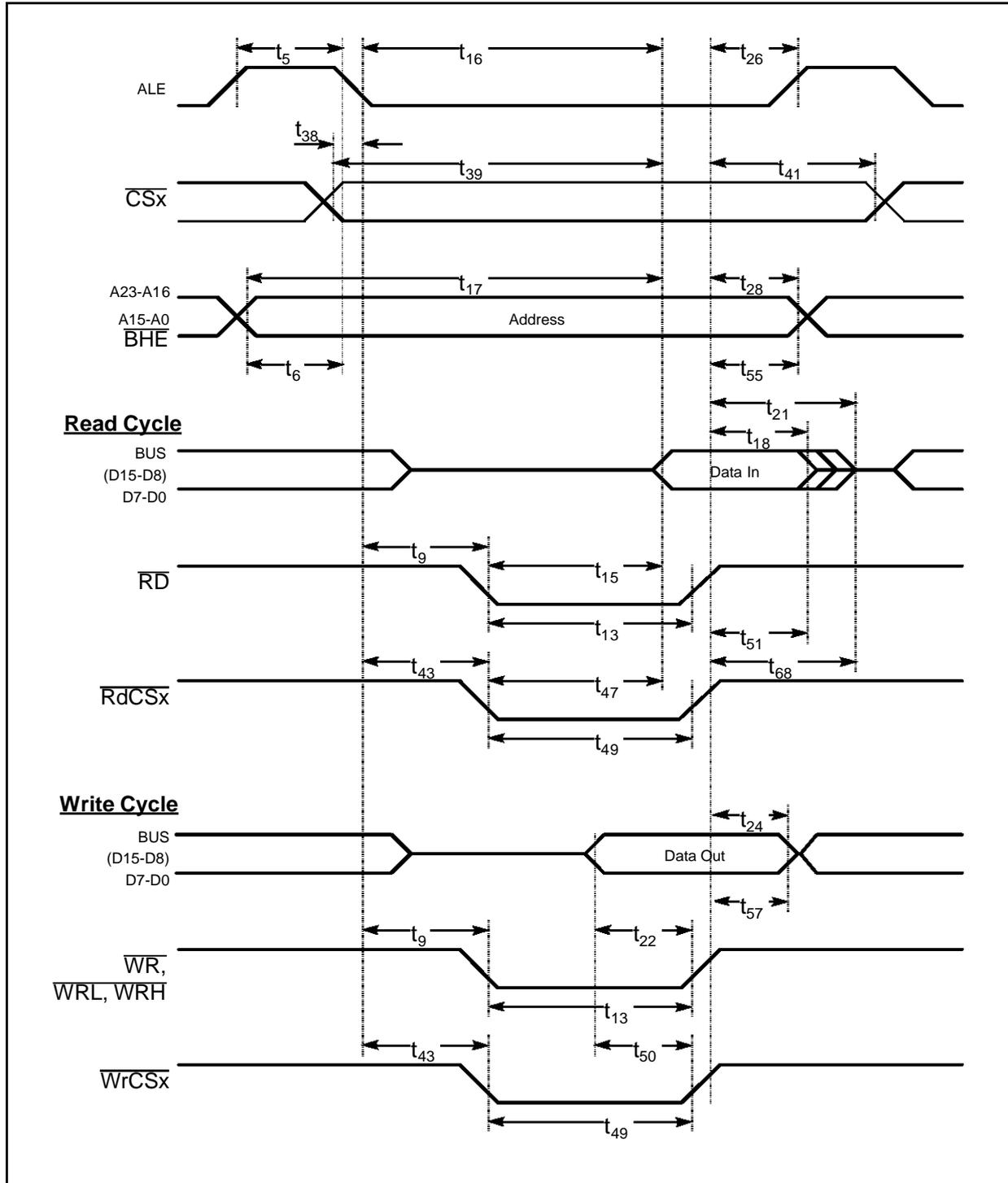


Figure 19. External Memory Cycle: Demultiplexed Bus, No Read/Write Delay, Extended ALE



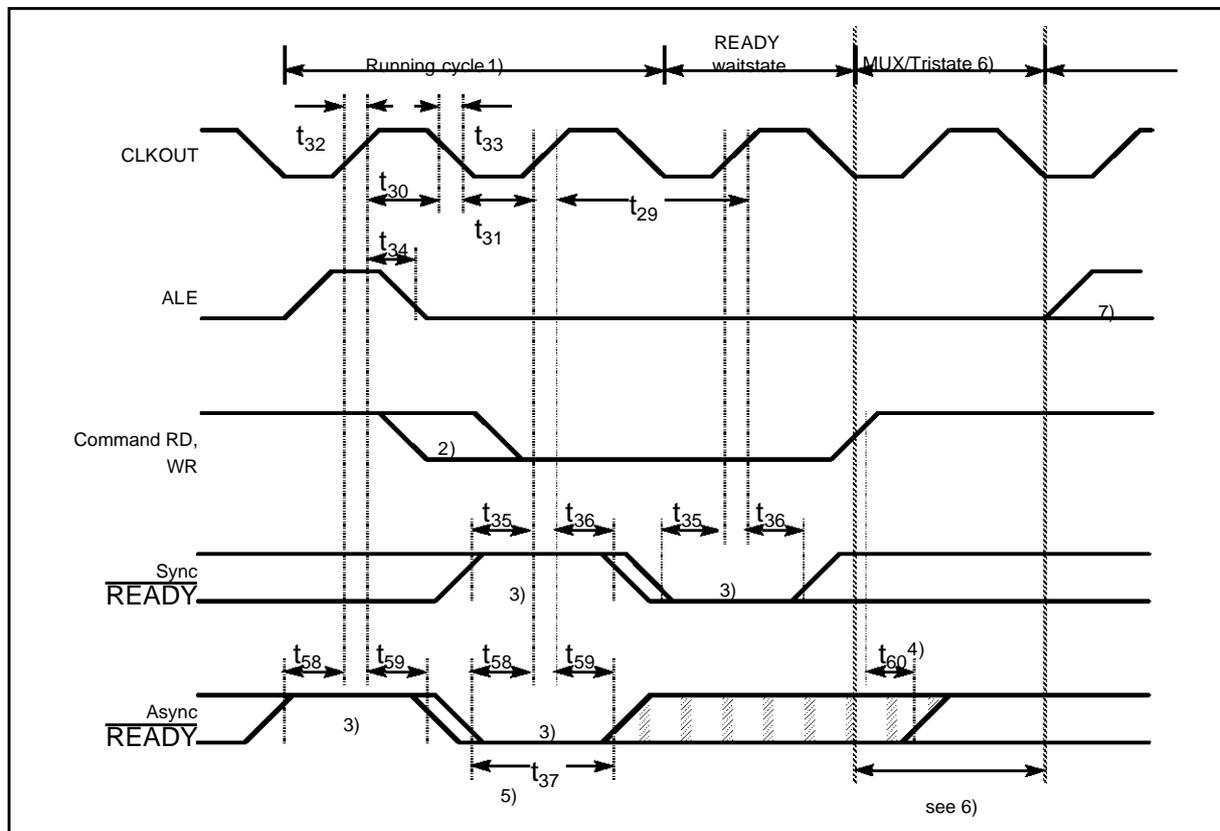
## AC Characteristics (Cont'd)

CLKOUT and  $\overline{\text{READY}}$  $V_{DD} = 5 \text{ V} \pm 10\%; V_{SS} = 0 \text{ V}$  $T_A = 0 \text{ to } +70 \text{ }^\circ\text{C}$  $C_L$  (for PORT0, PORT1, Port 4, ALE,  $\overline{\text{RD}}$ ,  $\overline{\text{WR}}$ ,  $\overline{\text{BHE}}$ , CLKOUT) = 100 pF $C_L$  (for Port 6,  $\overline{\text{CS}}$ ) = 100 pF

Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
CLKOUT cycle time	$t_{29}$ CC	50	50	2TCL	2TCL	ns
CLKOUT high time	$t_{30}$ CC	20	–	TCL – 5	–	ns
CLKOUT low time	$t_{31}$ CC	15	–	TCL – 10	–	ns
CLKOUT rise time	$t_{32}$ CC	–	5	–	5	ns
CLKOUT fall time	$t_{33}$ CC	–	5	–	5	ns
CLKOUT rising edge to ALE falling edge	$t_{34}$ CC	$0 + t_A$	$10 + t_A$	$0 + t_A$	$10 + t_A$	ns
Synchronous $\overline{\text{READY}}$ setup time to CLKOUT	$t_{35}$ SR	10	–	10	–	ns
Synchronous $\overline{\text{READY}}$ hold time after CLKOUT	$t_{36}$ SR	0	–	0	–	ns
Asynchronous $\overline{\text{READY}}$ low time	$t_{37}$ SR	65	–	2TCL + 15	–	ns
Asynchronous $\overline{\text{READY}}$ setup time <sup>1)</sup>	$t_{58}$ SR	15	–	15	–	ns
Asynchronous $\overline{\text{READY}}$ hold time <sup>1)</sup>	$t_{59}$ SR	0	–	0	–	ns
Async. $\overline{\text{READY}}$ hold time after $\overline{\text{RD}}$ , $\overline{\text{WR}}$ high (Demul- tiplexed Bus) <sup>2)</sup>	$t_{60}$ SR	0	$0 + 2t_{A2} + t_F$	0	$\text{TCL} - 25 + 2t_{A2} + t_F$	ns

**Notes:**

- 1) These timings are given for test purposes only, in order to assure recognition at a specific clock edge.
- 2) Demultiplexed bus is the worst case. For multiplexed bus 2TCL are to be added to the maximum values. This adds even more time for deactivating  $\overline{\text{READY}}$ .  
The  $2t_A$  refer to the next following bus cycle.

Figure 20. CLKOUT and  $\overline{\text{READY}}$ **Notes:**

- 1) Cycle as programmed, including MCTC waitstates (Example shows 0 MCTC WS).
- 2) The leading edge of the respective command depends on RW-delay.
- 3)  $\overline{\text{READY}}$  sampled HIGH at this sampling point generates a READY controlled waitstate,  $\overline{\text{READY}}$  sampled LOW at this sampling point terminates the currently running bus cycle.
- 4)  $\overline{\text{READY}}$  may be deactivated in response to the trailing (rising) edge of the corresponding command ( $\overline{\text{RD}}$  or  $\overline{\text{WR}}$ ).
- 5) If the Asynchronous  $\overline{\text{READY}}$  signal does not fulfill the indicated setup and hold times with respect to CLKOUT (e.g. because CLKOUT is not enabled), it must fulfill  $t_{37}$  in order to be safely synchronized. This is guaranteed, if READY is removed in response to the command (see Note 4)).
- 6) Multiplexed bus modes have a MUX waitstate added after a bus cycle, and an additional MTTC waitstate may be inserted here.  
For a multiplexed bus with MTTC waitstate this delay is 2 CLKOUT cycles, for a demultiplexed bus without MTTC waitstate this delay is zero.
- 7) The next external bus cycle may start here.

**13.5.2.2 External Bus Arbitration**

$V_{DD} = 5\text{ V} \pm 10\%$ ;  $V_{SS} = 0\text{ V}$

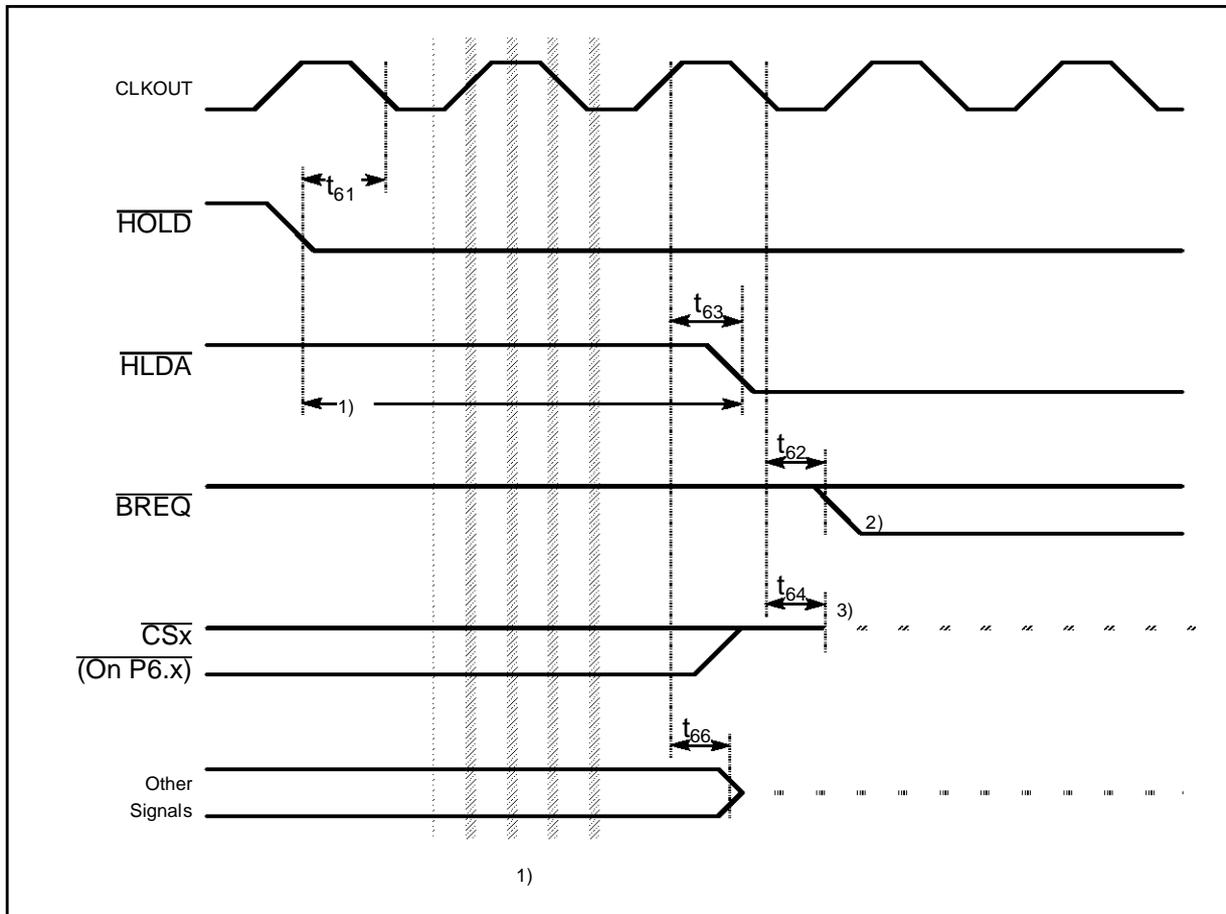
$T_A = 0\text{ to }+70\text{ }^\circ\text{C}$

$C_L$  (for PORT0, PORT1, Port 4, ALE,  $\overline{\text{RD}}$ ,  $\overline{\text{WR}}$ ,  $\overline{\text{BHE}}$ , CLKOUT) = 100 pF

$C_L$  (for Port 6,  $\overline{\text{CS}}$ ) = 100 pF

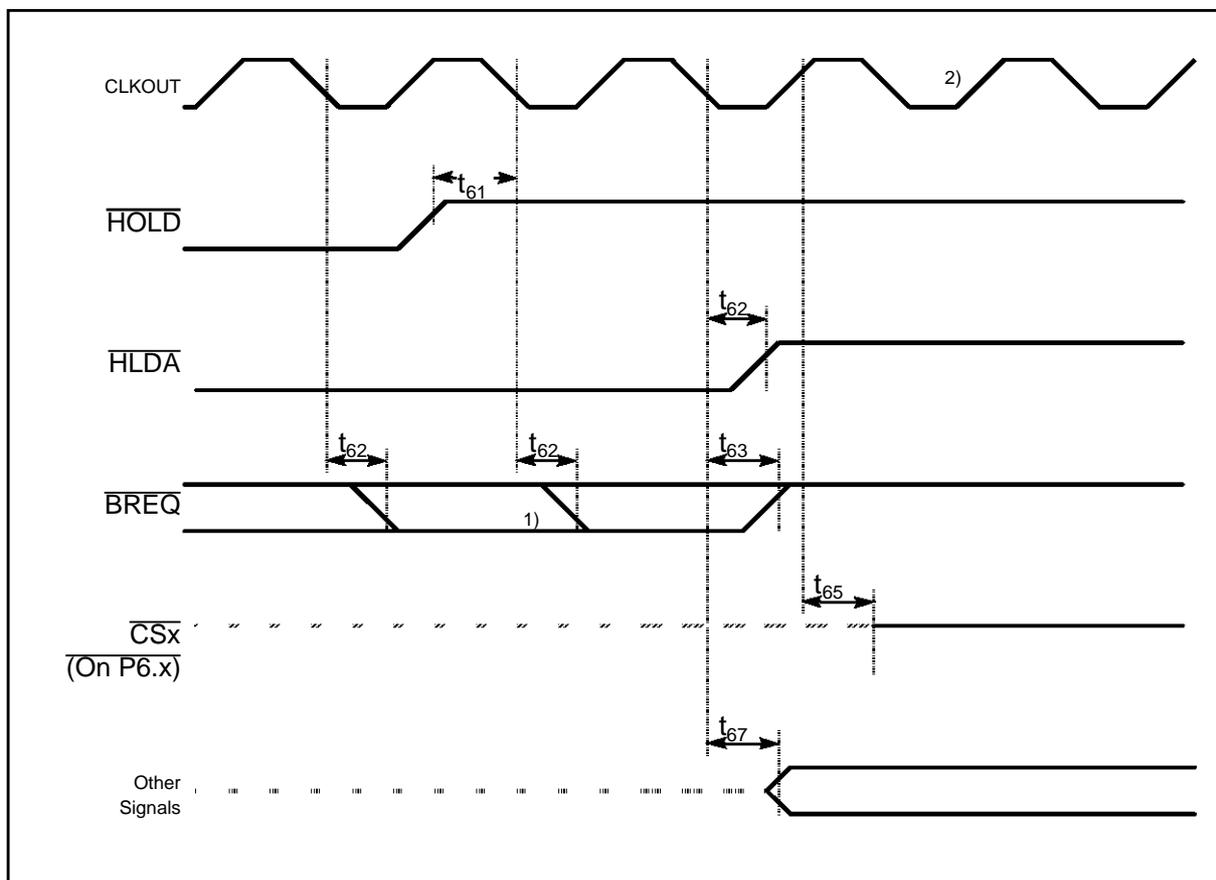
Parameter	Symbol	Max. CPU Clock = 20 MHz		Variable CPU Clock 1/2TCL = 1 to 20 MHz		Unit
		min.	max.	min.	max.	
HOLD input setup time to CLKOUT	t <sub>61</sub> SR	20	–	20	–	ns
CLKOUT to $\overline{\text{HLD}}\overline{\text{A}}$ high or $\overline{\text{BREQ}}$ low delay	t <sub>62</sub> CC	–	20	–	20	ns
CLKOUT to $\overline{\text{HLD}}\overline{\text{A}}$ low or $\overline{\text{BREQ}}$ high delay	t <sub>63</sub> CC	–	20	–	20	ns
$\overline{\text{CS}}\overline{\text{x}}$ release	t <sub>64</sub> CC	–	20	–	20	ns
$\overline{\text{CS}}\overline{\text{x}}$ drive	t <sub>65</sub> CC	-5	25	-5	25	ns
Other signals release	t <sub>66</sub> CC	–	20	–	20	ns
Other signals drive	t <sub>67</sub> CC	-5	25	-5	25	ns

Figure 21. External Bus Arbitration, Releasing the Bus

**Notes:**

- 1) The ST10R165 will complete the currently running bus cycle before granting bus access.
- 2) This is the first possibility for  $\overline{\text{BREQ}}$  to get active.
- 3) The  $\overline{\text{CS}}$  outputs will be resistive high (pullup) after  $t_{64}$ .

Figure 22. External Bus Arbitration, (Regaining the Bus)



**Notes:**

- 1) This is the last chance for  $\overline{BREQ}$  to trigger the indicated regain-sequence. Even if  $\overline{BREQ}$  is activated earlier, the regain-sequence is initiated by  $\overline{HOLD}$  going high. Please note that  $\overline{HOLD}$  may also be deactivated without the ST10R165 requesting the bus.
- 2) The next ST10R165 driven bus cycle may start here.

14 PACKAGE MECHANICAL DATA

Figure 23. Package Outline PQFP100 (14 x 20 mm)

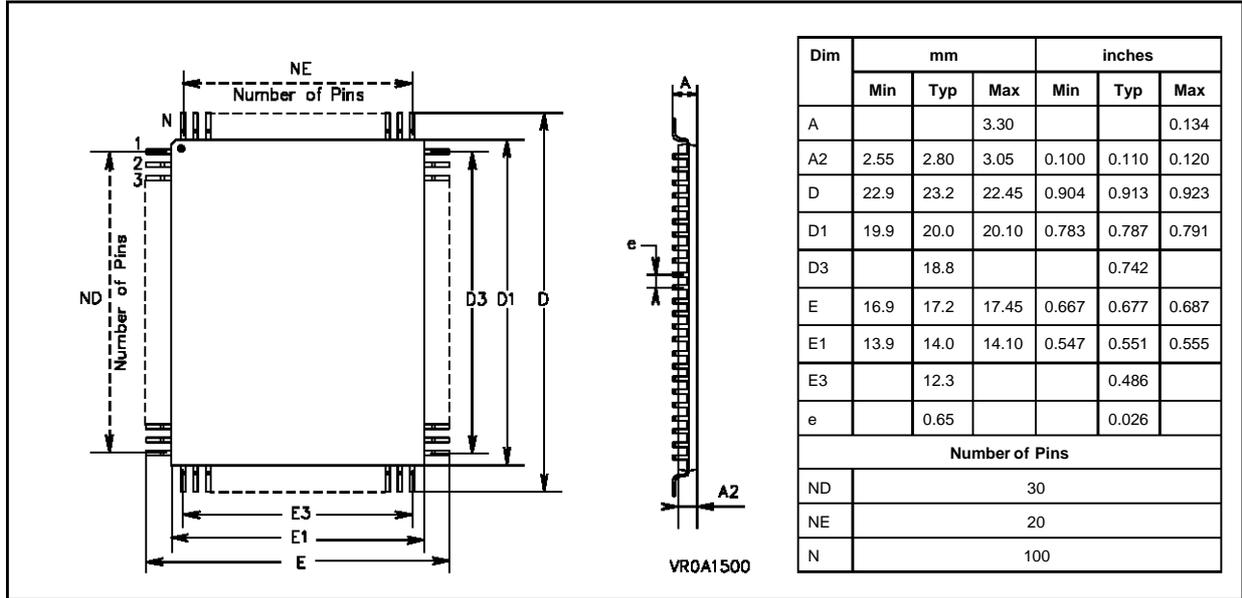
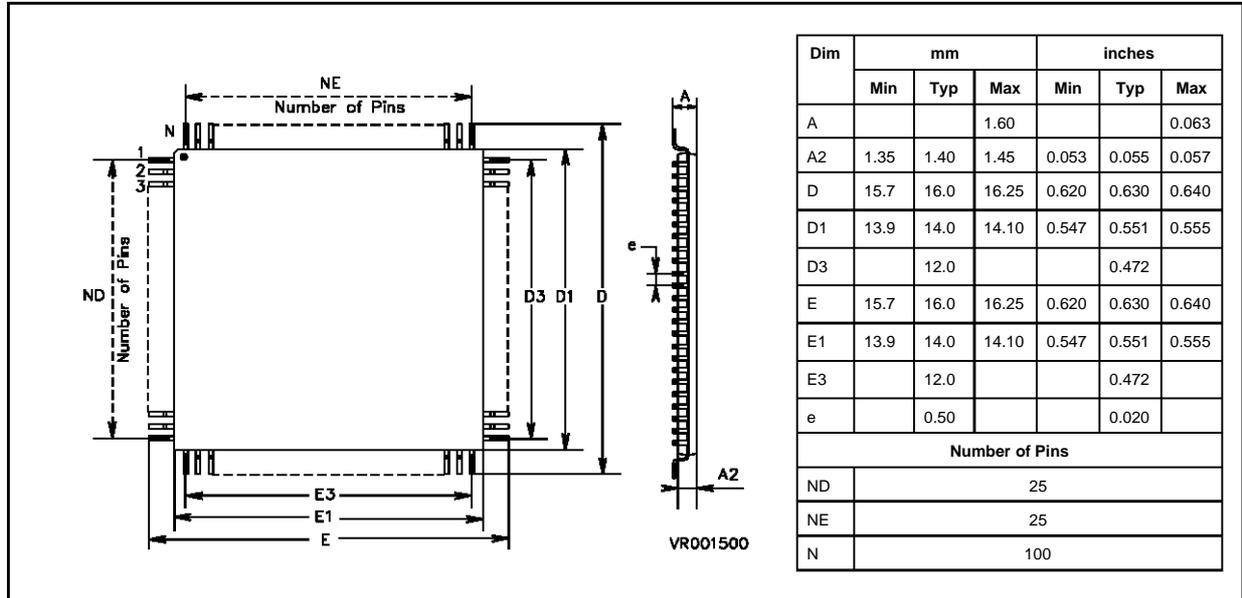


Figure 24. Package Outline TQFP100 (14 x 14 mm)



**15 ORDERING INFORMATION**

<b>Salestype</b>	<b>Temperature range</b>	<b>Package</b>
ST10R165BQ1	0°C to 70°C	PQFP100(14x 20)
ST10R165BQ6	- 40°C to 85°C	PQFP100(14x 20)
ST10R165BT1	0°C to 70°C	TQFP100(14x 14)
ST10R165BT6	- 40°C to 85°C	TQFP100(14x 14)

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